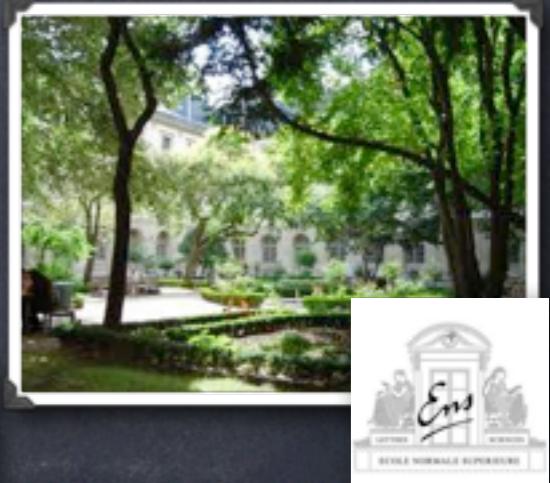


Non local causality analysis in rule-based models

Jean Krivine
CNRS & Univ. Paris Diderot

“Programs as models” project

- Import tools from theoretical computer science in order to describe and analyze biological systems
- Develop new technique to cope with the complexity of the cell



Theoretical computer science

Abstract interpretation

Concurrency theory

Context free languages

Intrinsic Information Carriers in Combinatorial Dynamical Systems

Russ Harmer,^{1,2} Vincent Danos,³ Jérôme Feret,⁴ Jean Krivine,² and Walter Fontana¹

Internal coarse-graining of molecular systems

Jérôme Feret ^{*}, Vincent Danos [†], Jean Krivine ^{*}, Russ Harmer [‡], and Walter Fontana ^{*}

^{*}Harvard Medical School, Boston, USA, [†]University of Edinburgh, Edinburgh, United Kingdom, and [‡]CNRS & Paris Diderot, Paris, France

Submitted to Proceedings of the National Academy of Sciences of the United States of America

Rule-based modelling of cellular signalling

Vincent Danos^{1,3,4}, Jérôme Feret², Walter Fontana³, Russell Harmer^{3,4}, and Jean Krivine⁵

Scalable simulation of cellular signaling networks

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KaSim 3.0

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Today's talk!

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KaSim 3.0

Causality analysis

The easy life of a computer scientist

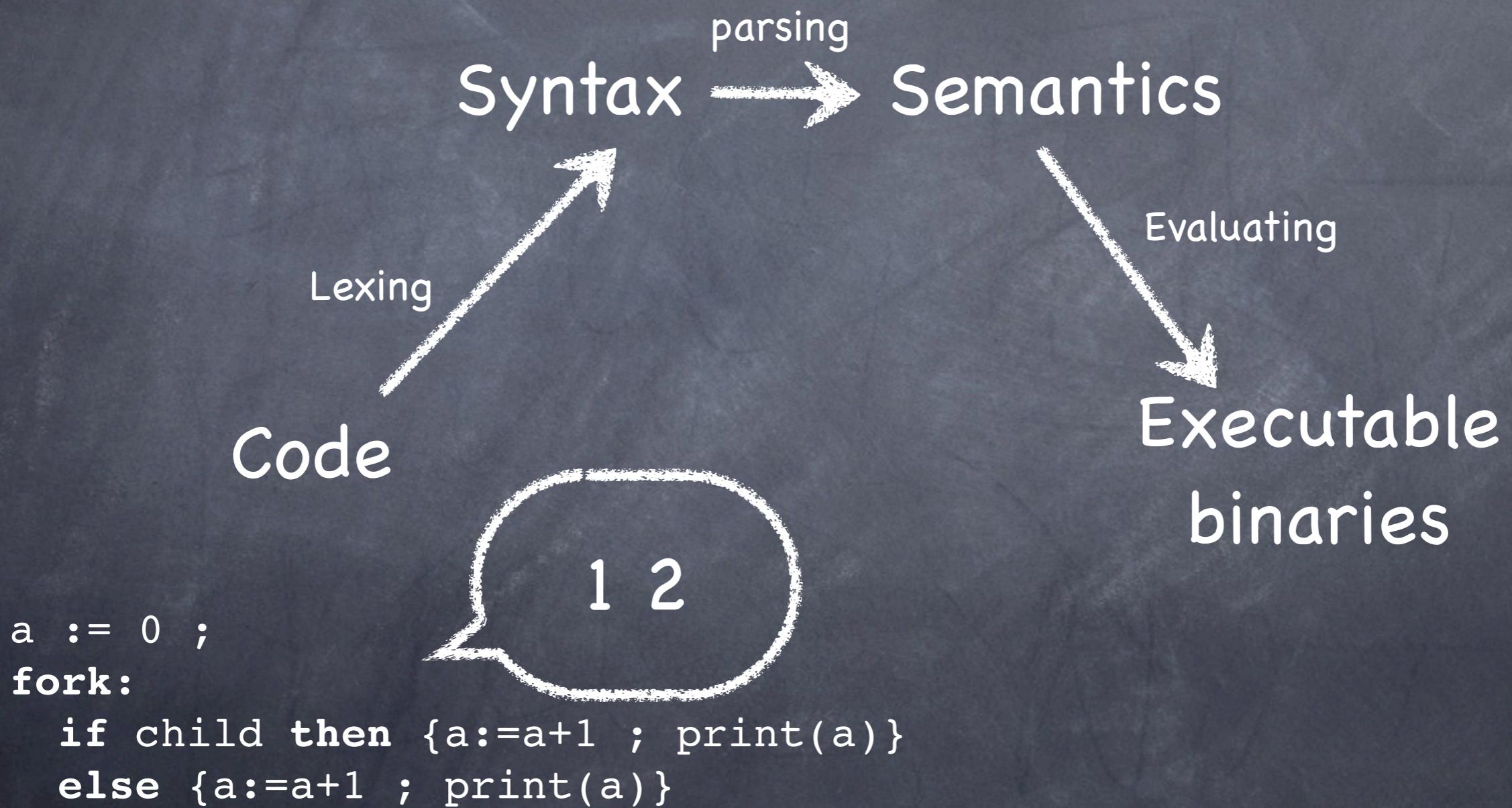
```
a := 0 ;
fork:
  if child then {a:=a+1 ; print(a)}
  else {a:=a+1 ; print(a)}
```

The easy life of a computer scientist

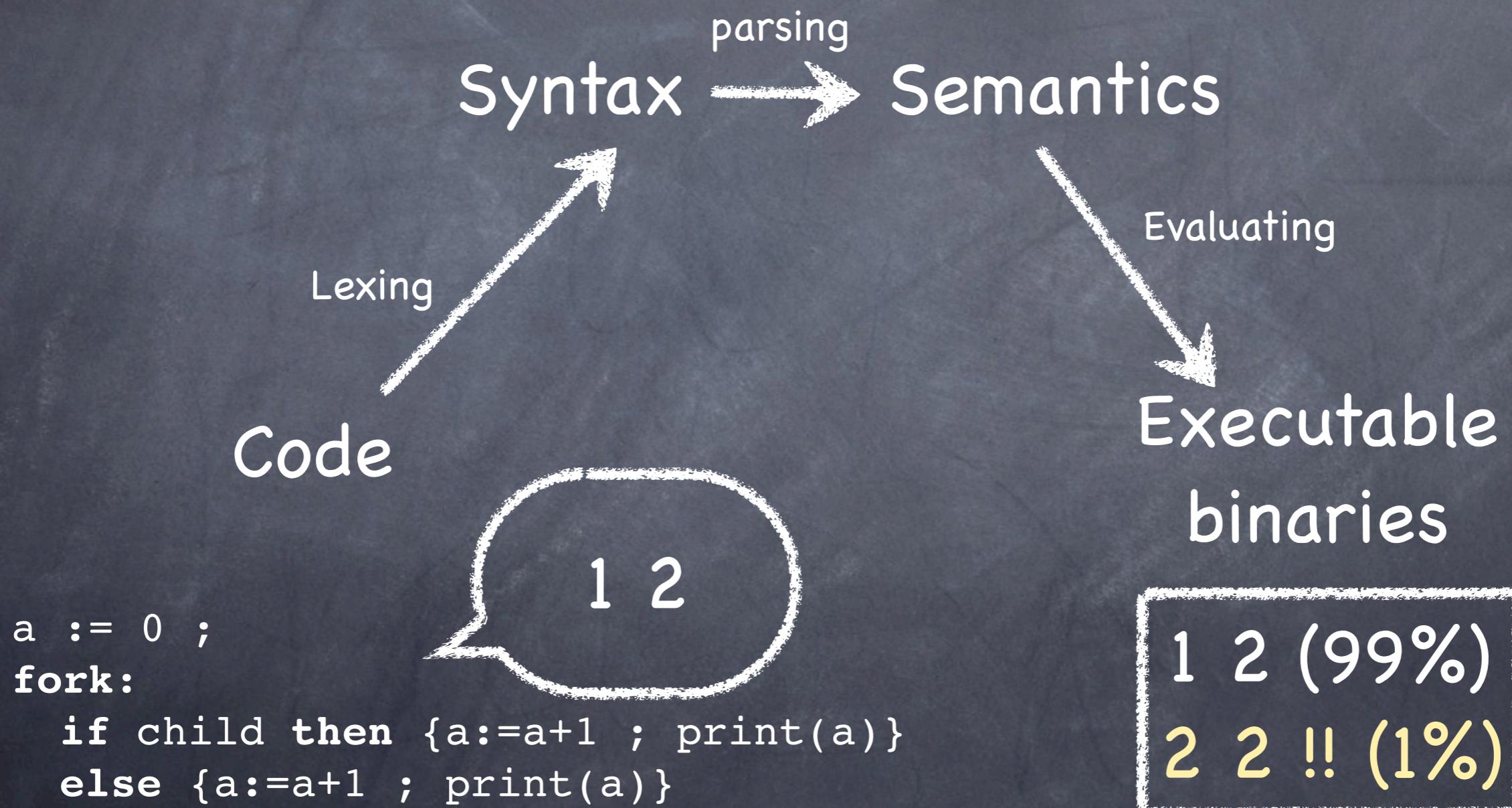
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```



The easy life of a computer scientist



The easy life of a computer scientist



Debugging

T0: set a:=0

T0: spawn T1

T0: set a:=a+1

T1: set a:=a+1

T1: print (a)

T1: terminate

T0: print (a)

observation: 2 2



Interleaving semantics

Debugging

T0: set a:=0

T0: spawn T1

T0: set a:=a+1

T1: set a:=a+1

T1: print (a)

T1: terminate

T0: print (a)

observation: 2 2



Interleaving semantics

Non-Interleaving semantics

Debugging

T0: set a:=0

T0: spawn T1

T0: set a:=a+1

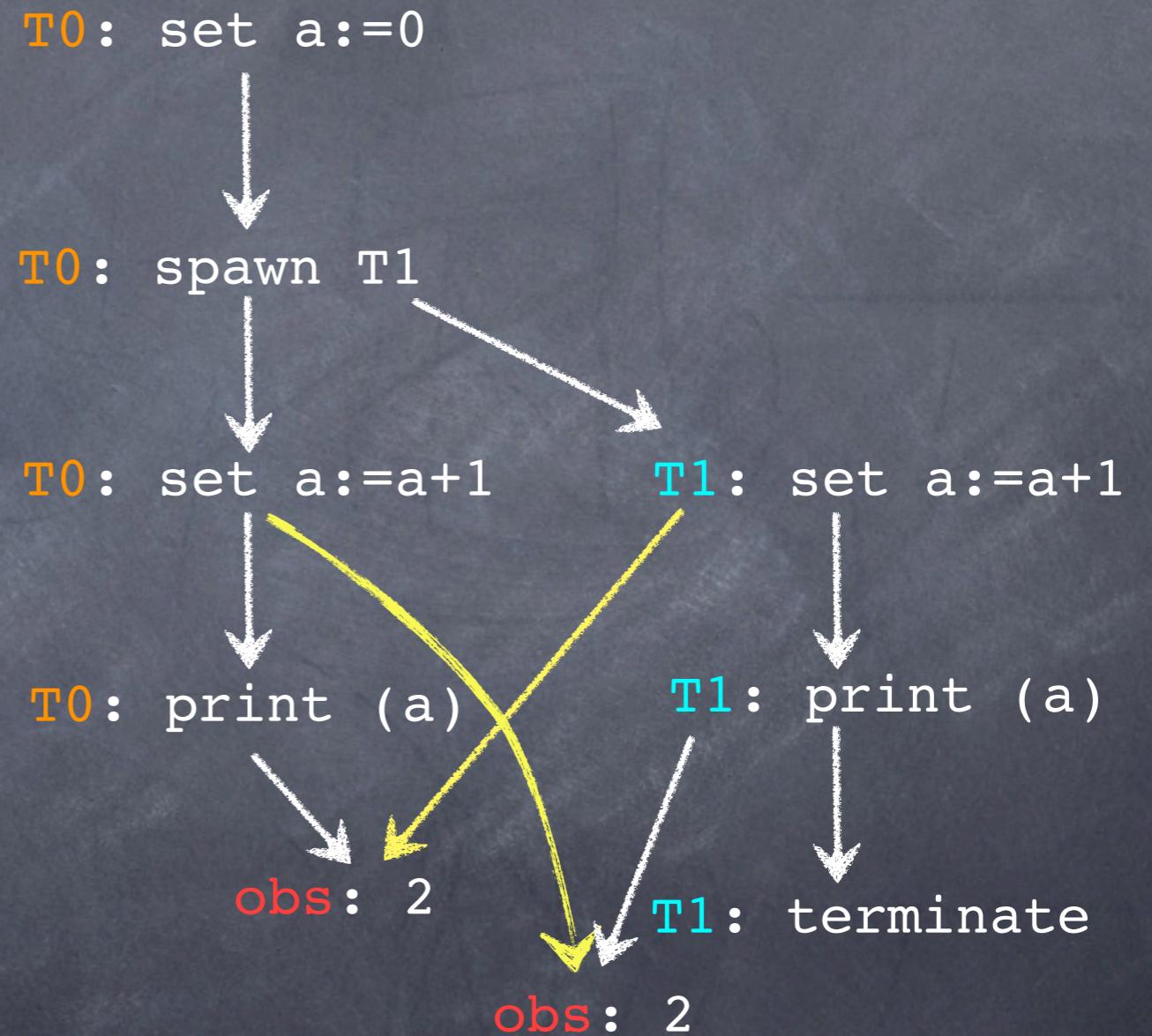
T1: set a:=a+1

T1: print (a)

T1: terminate

T0: print (a)

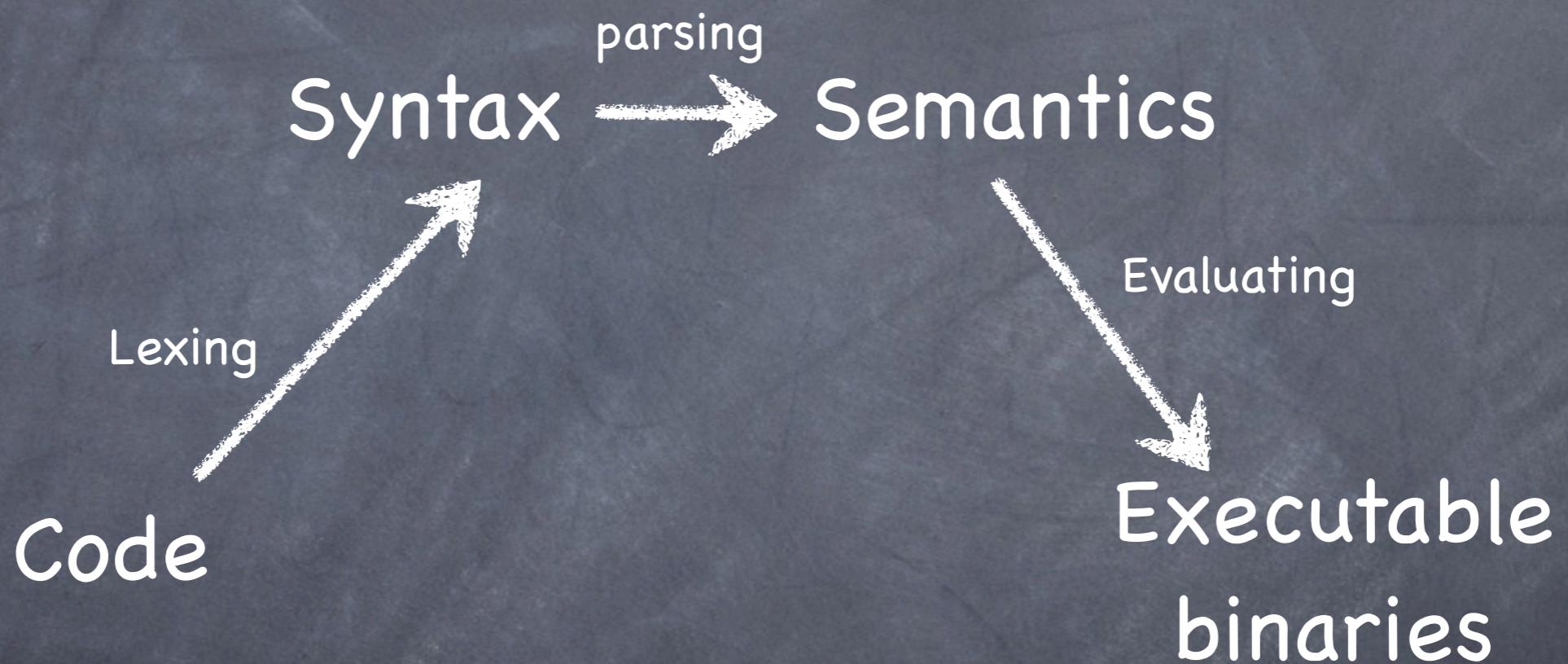
observation: 2 2



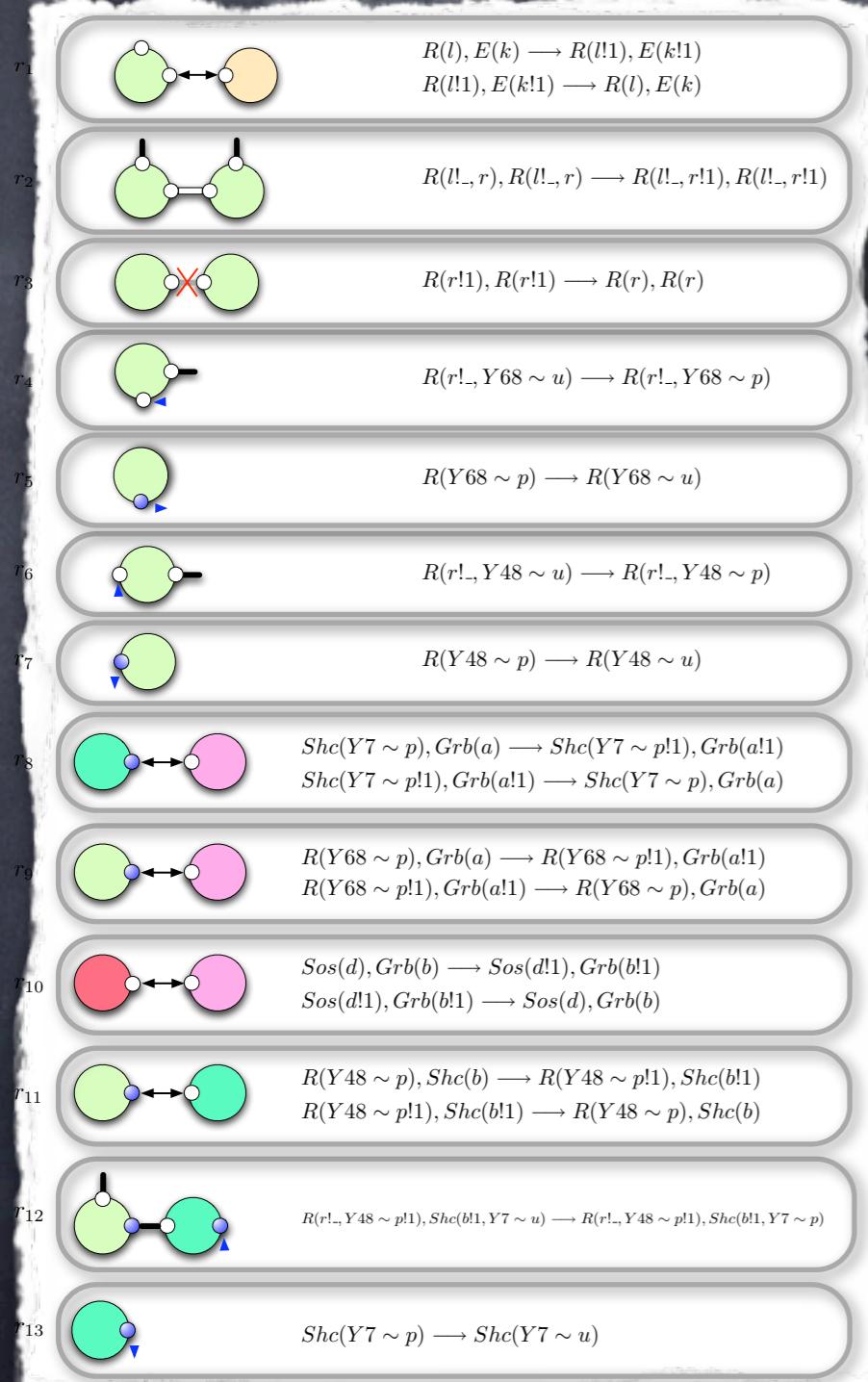
Interleaving semantics

Non-Interleaving semantics

The difficult life of a modeler



The difficult life of a modeler

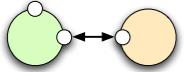
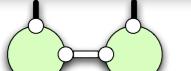
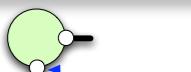
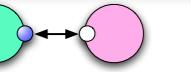
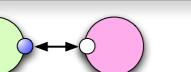
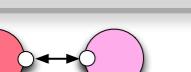


yntax $\xrightarrow{\text{parsing}}$ Semantics

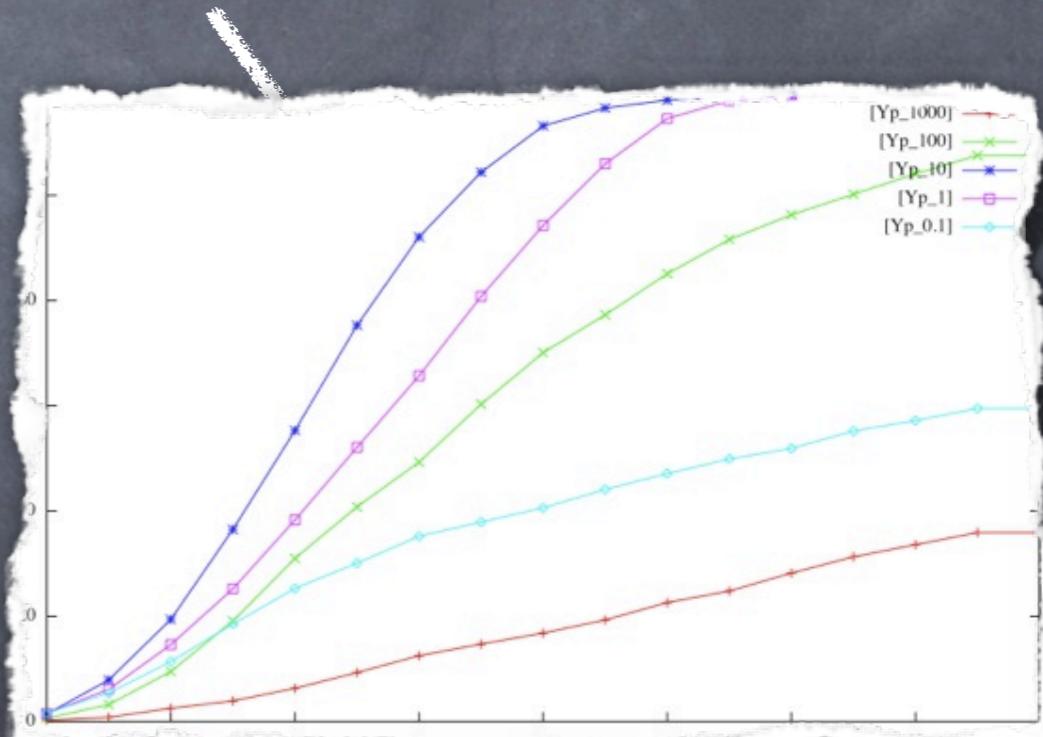
Evaluating

Executable
binaries

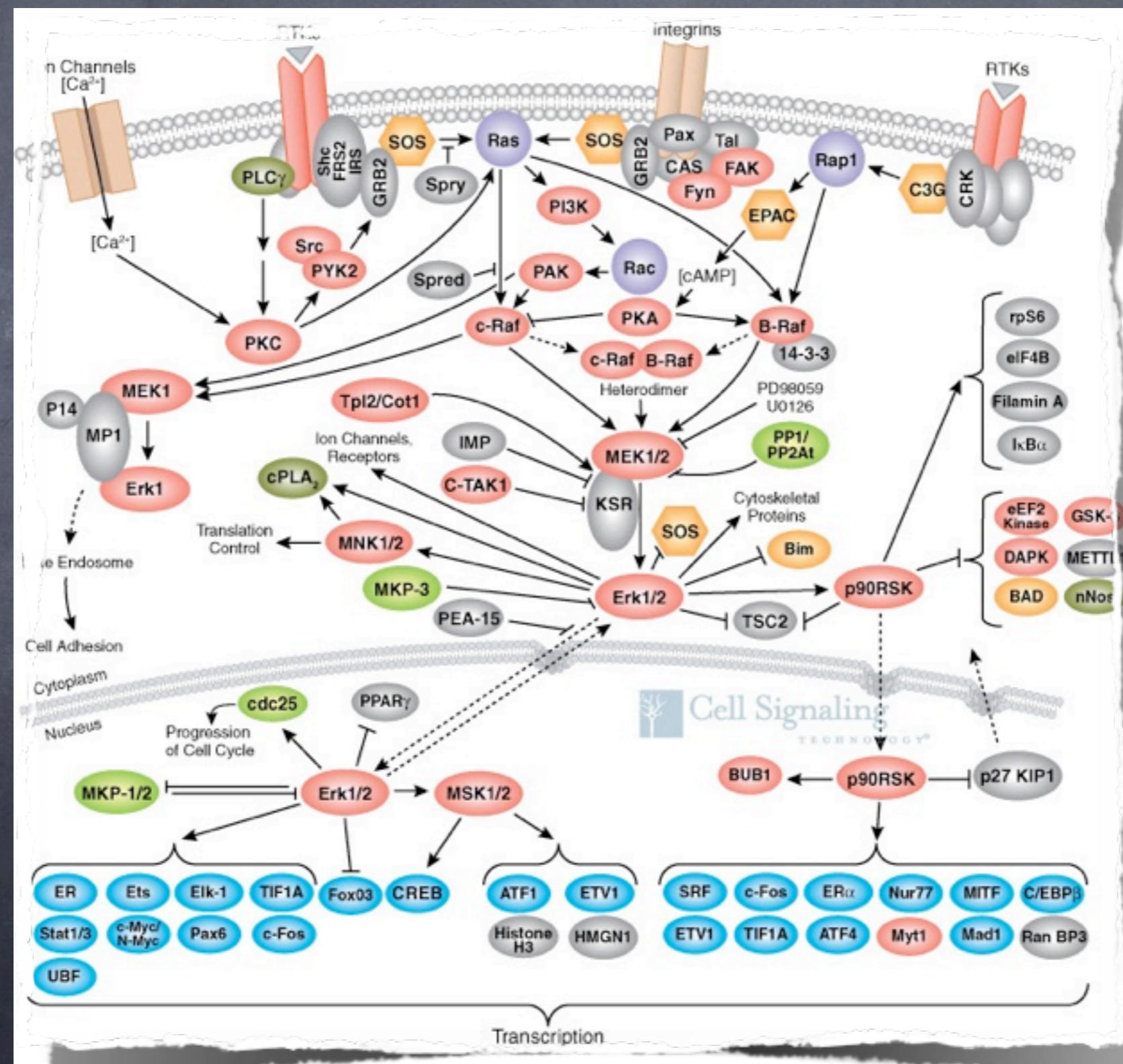
The difficult life of a modeler

r_1		$R(l), E(k) \longrightarrow R(l!1), E(k!1)$ $R(l!1), E(k!1) \longrightarrow R(l), E(k)$
r_2		$R(l!-, r), R(l!-, r) \longrightarrow R(l!-, r!1), R(l!-, r!1)$
r_3		$R(r!1), R(r!1) \longrightarrow R(r), R(r)$
r_4		$R(r!-, Y68 \sim u) \longrightarrow R(r!-, Y68 \sim p)$
r_5		$R(Y68 \sim p) \longrightarrow R(Y68 \sim u)$
r_6		$R(r!-, Y48 \sim u) \longrightarrow R(r!-, Y48 \sim p)$
r_7		$R(Y48 \sim p) \longrightarrow R(Y48 \sim u)$
r_8		$Shc(Y7 \sim p), Grb(a) \longrightarrow Shc(Y7 \sim p!1), Grb(a!1)$ $Shc(Y7 \sim p!1), Grb(a!1) \longrightarrow Shc(Y7 \sim p), Grb(a)$
r_9		$R(Y68 \sim p), Grb(a) \longrightarrow R(Y68 \sim p!1), Grb(a!1)$ $R(Y68 \sim p!1), Grb(a!1) \longrightarrow R(Y68 \sim p), Grb(a)$
r_{10}		$Sos(d), Grb(b) \longrightarrow Sos(d!1), Grb(b!1)$ $Sos(d!1), Grb(b!1) \longrightarrow Sos(d), Grb(b)$
r_{11}		$R(Y48 \sim p), Shc(b) \longrightarrow R(Y48 \sim p!1), Shc(b!1)$ $R(Y48 \sim p!1), Shc(b!1) \longrightarrow R(Y48 \sim p), Shc(b)$
r_{12}		$R(r!-, Y48 \sim p!1), Shc(b!1, Y7 \sim u) \longrightarrow R(r!-, Y48 \sim p!1), Shc(b!1, Y7 \sim p)$
r_{13}		$Shc(Y7 \sim p) \longrightarrow Shc(Y7 \sim u)$

syntax $\xrightarrow{\text{parsing}}$ Semantics

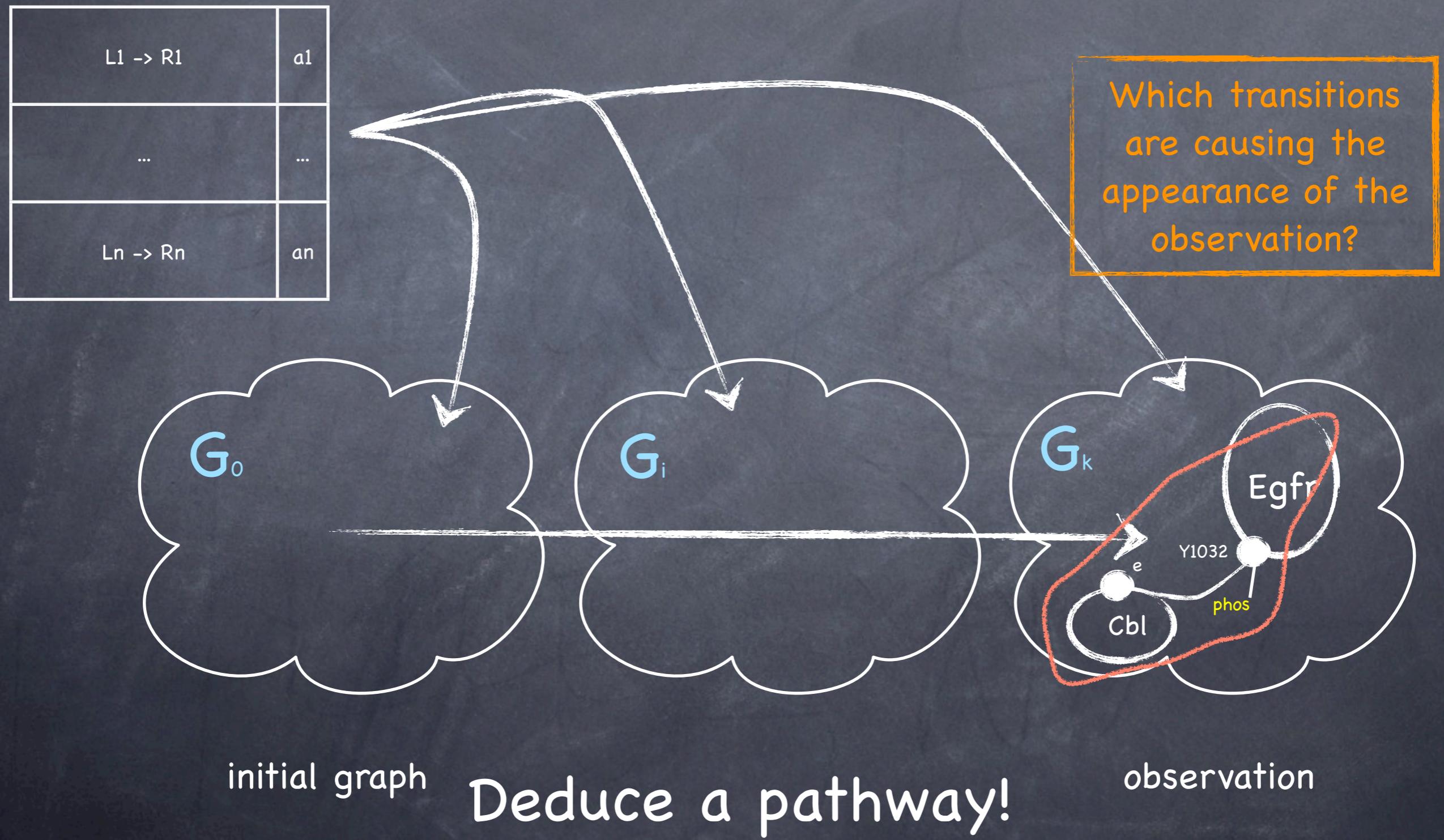


... and causality



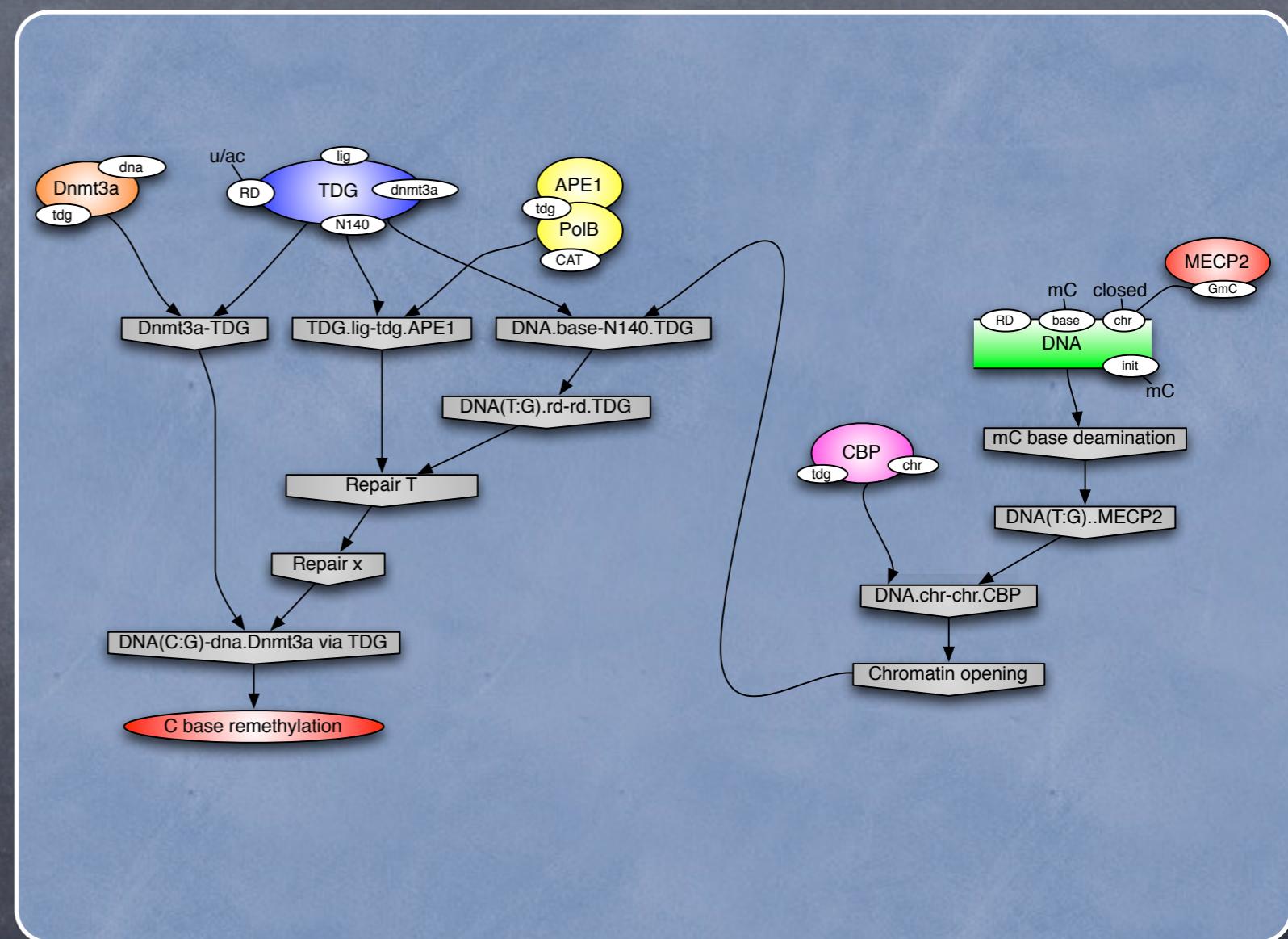
The question

The question



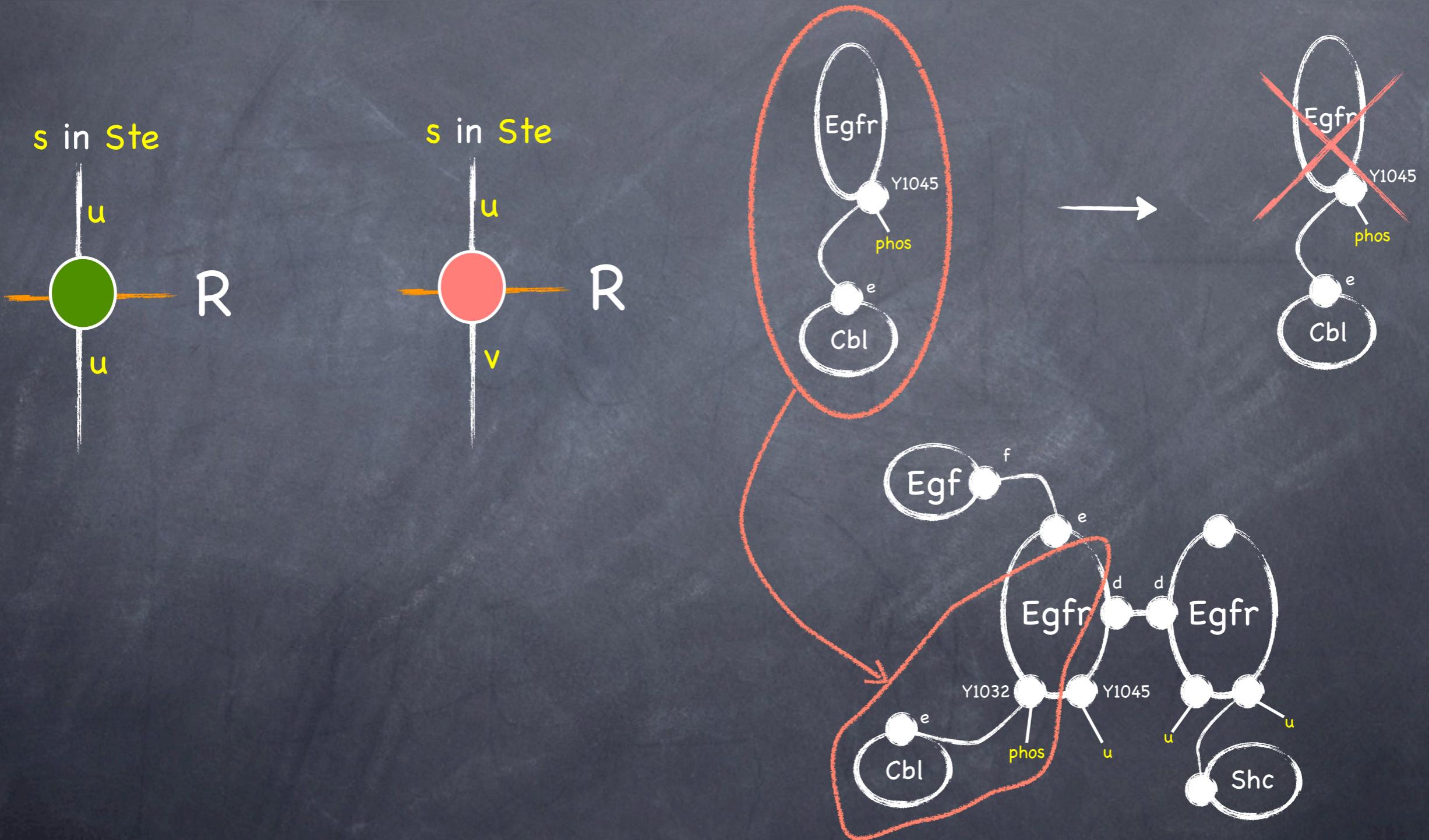
Remaining of the talk

- Classical causality analysis
- Notion of “knock-out” property
- Causal compression

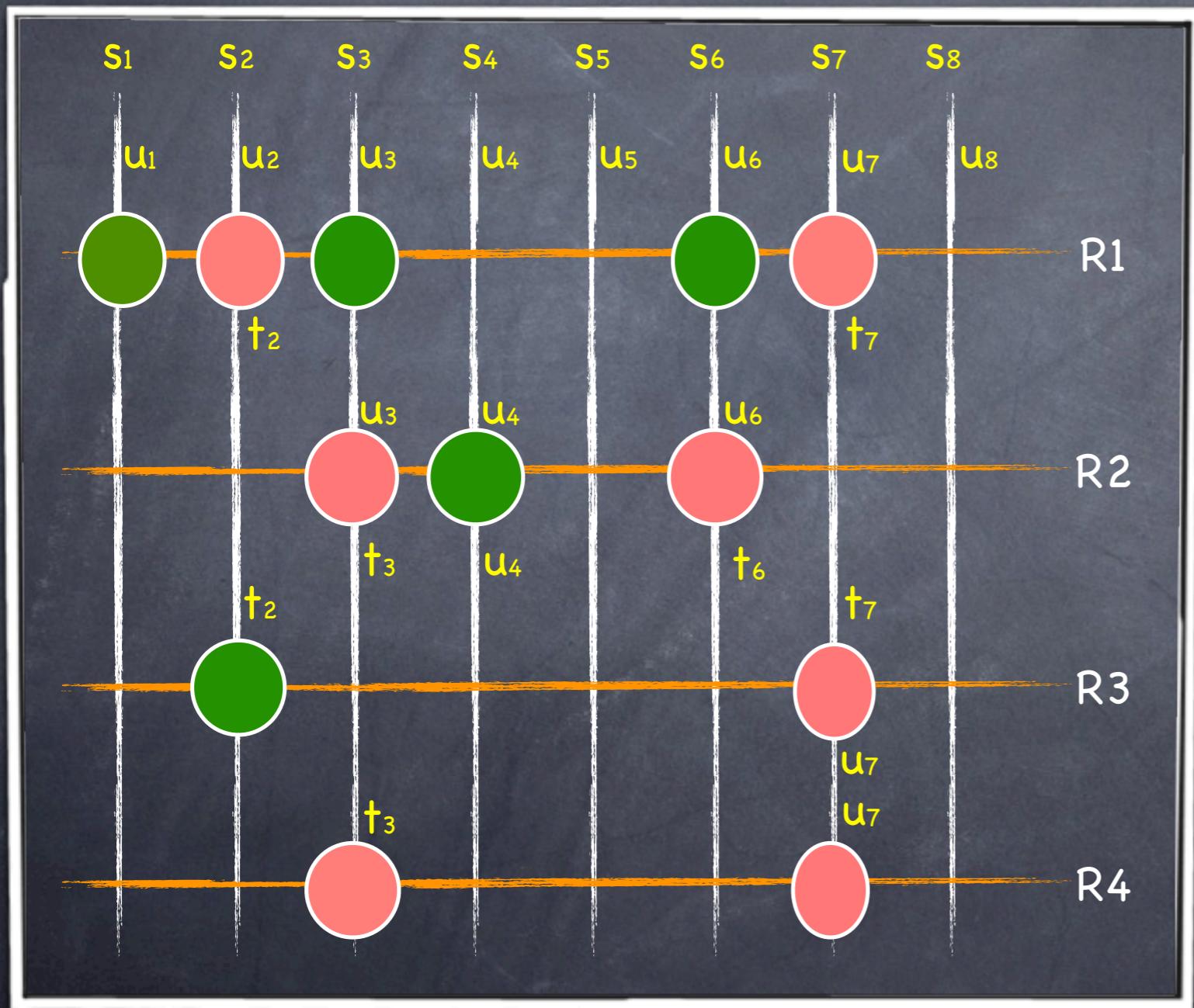


Classical causality analysis

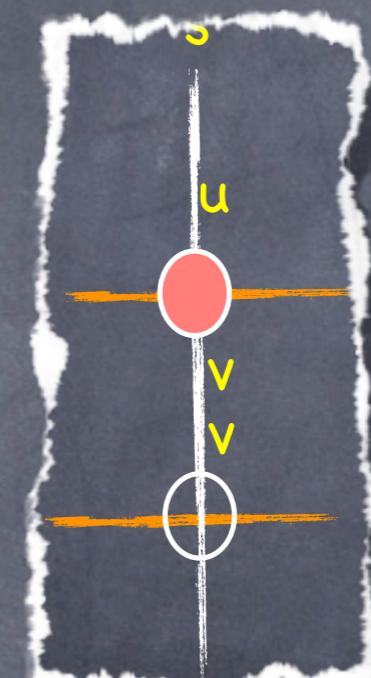
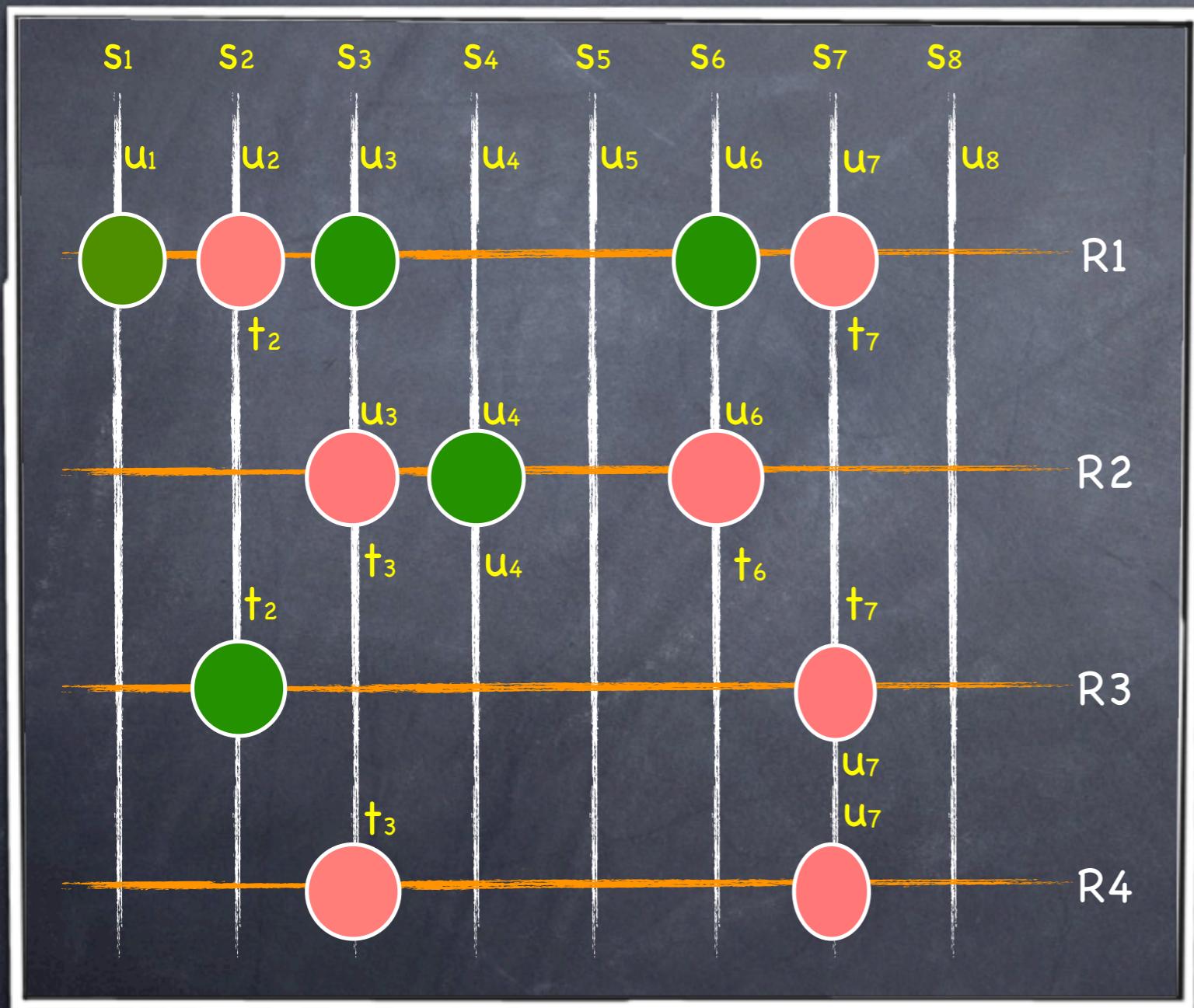
Tracking modifications



Dependencies

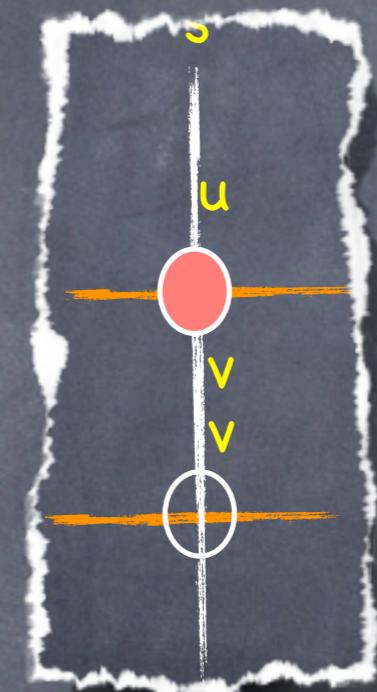
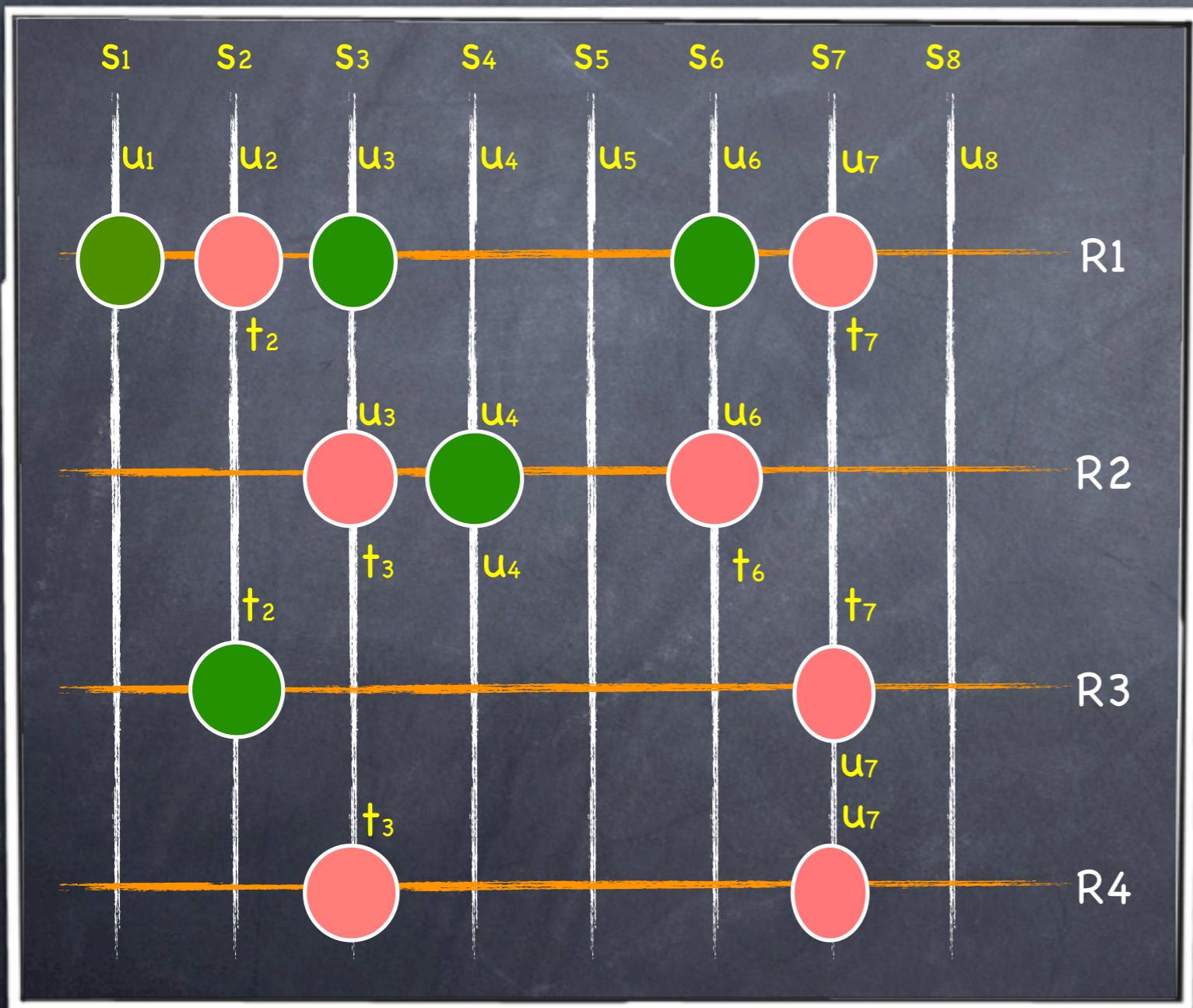


Dependencies

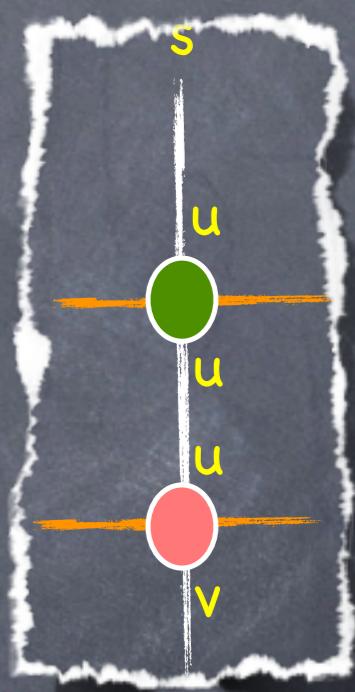


Causality

Dependencies

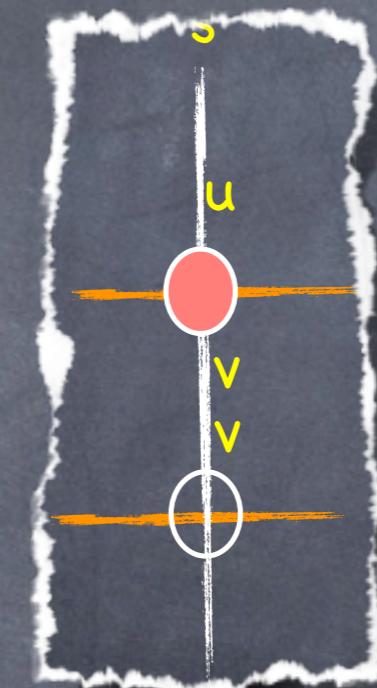
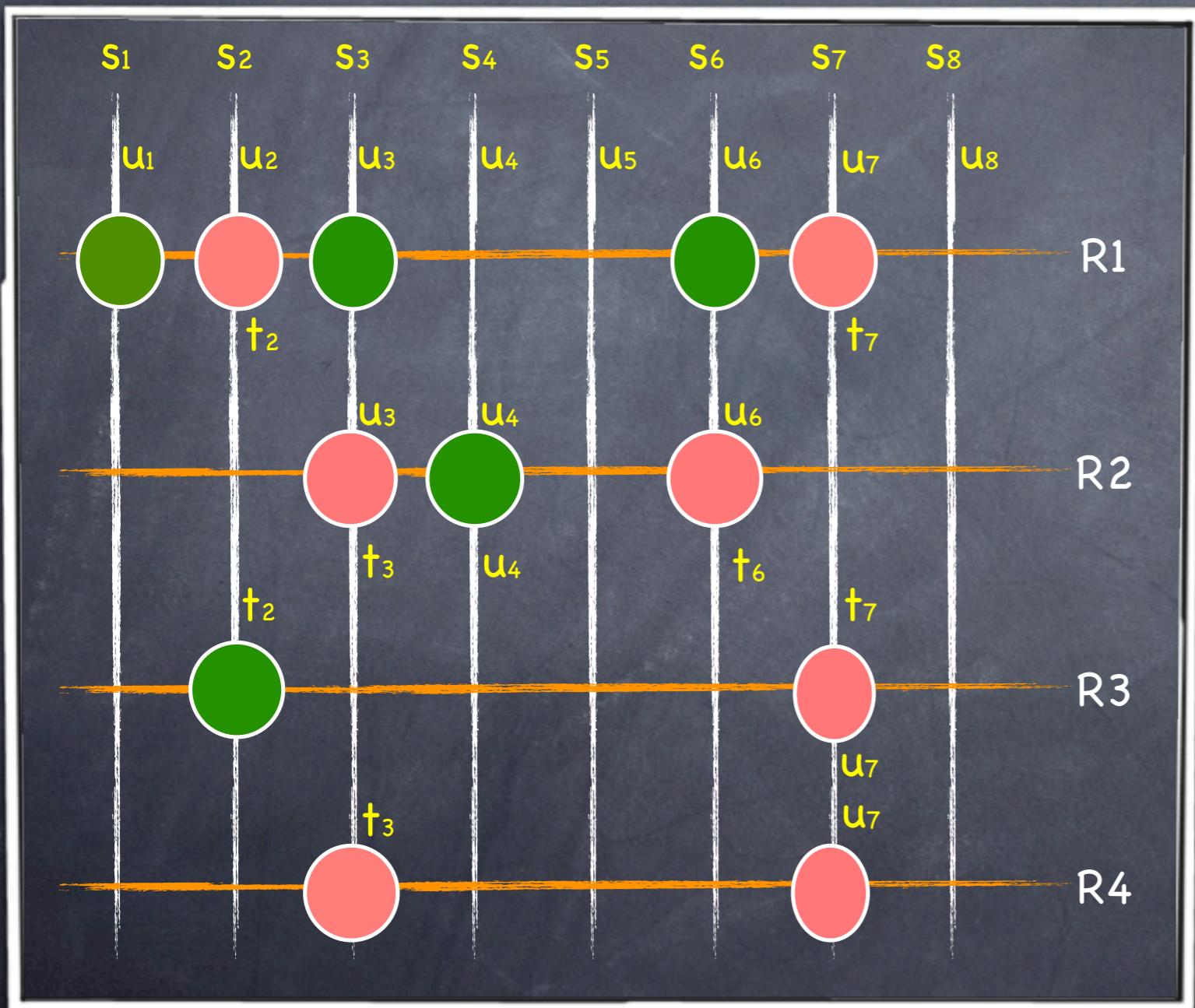


Causality

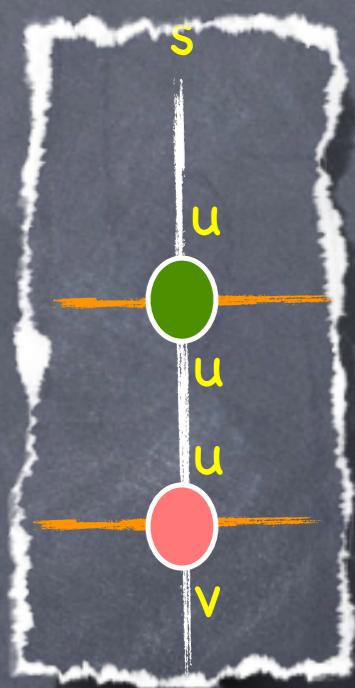


Precedence

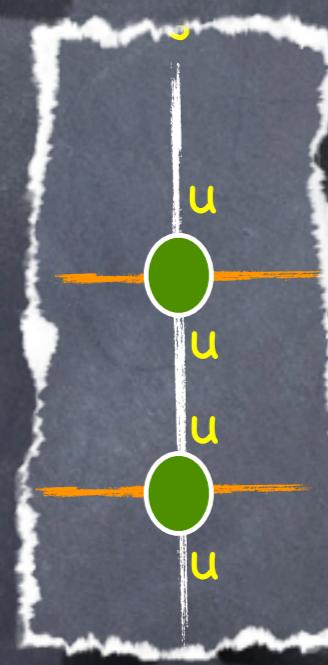
Dependencies



Causality

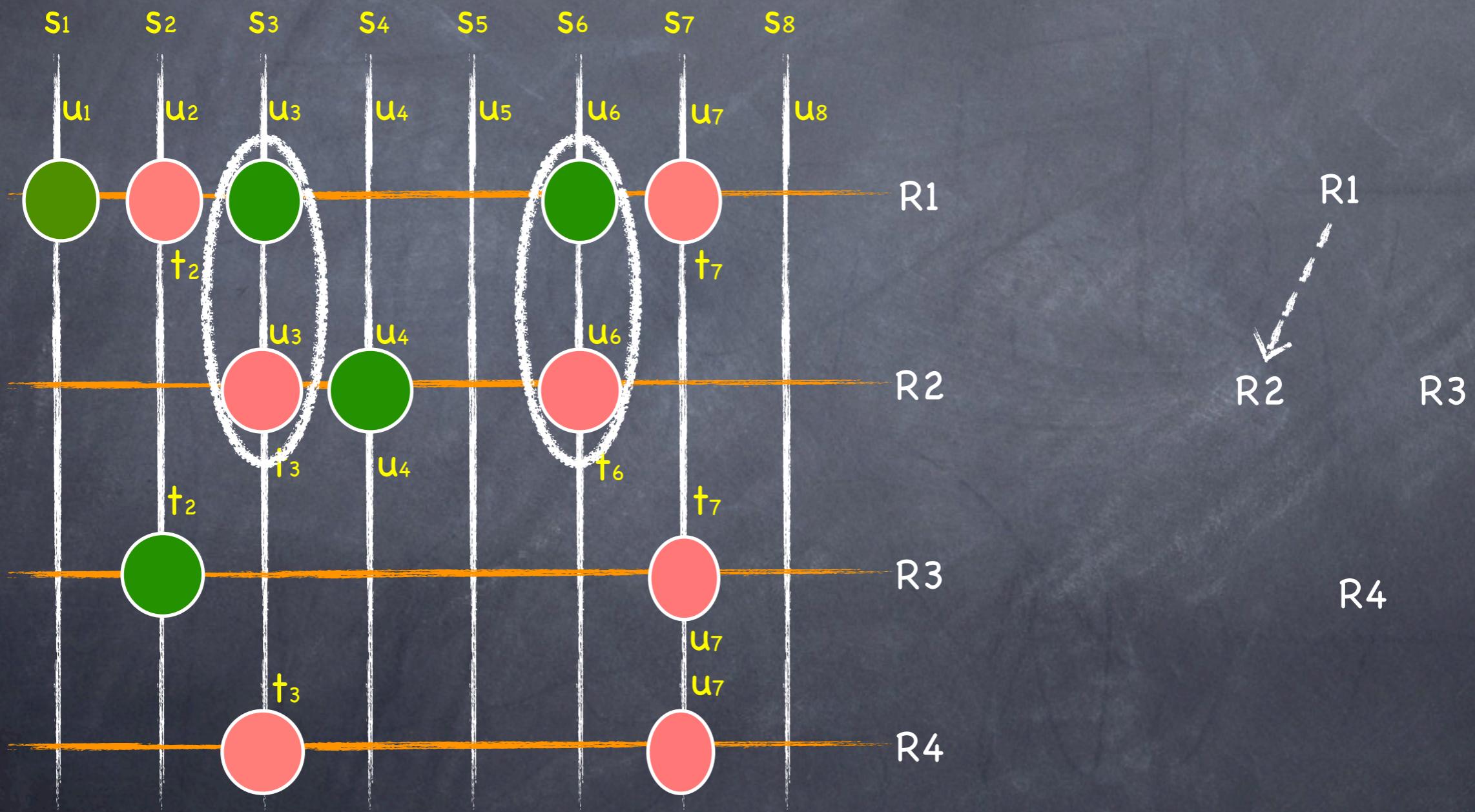


Precedence

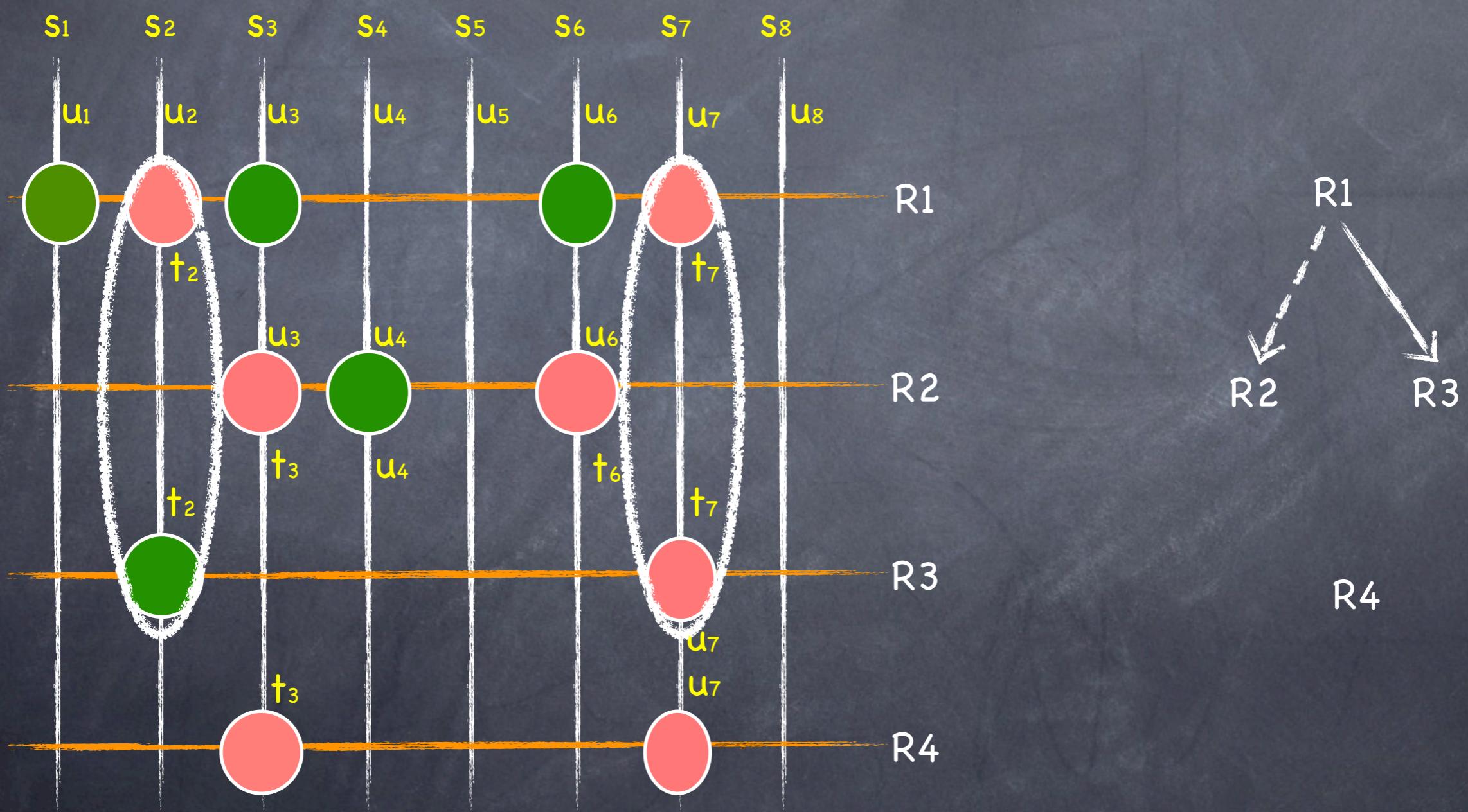


Concurrency

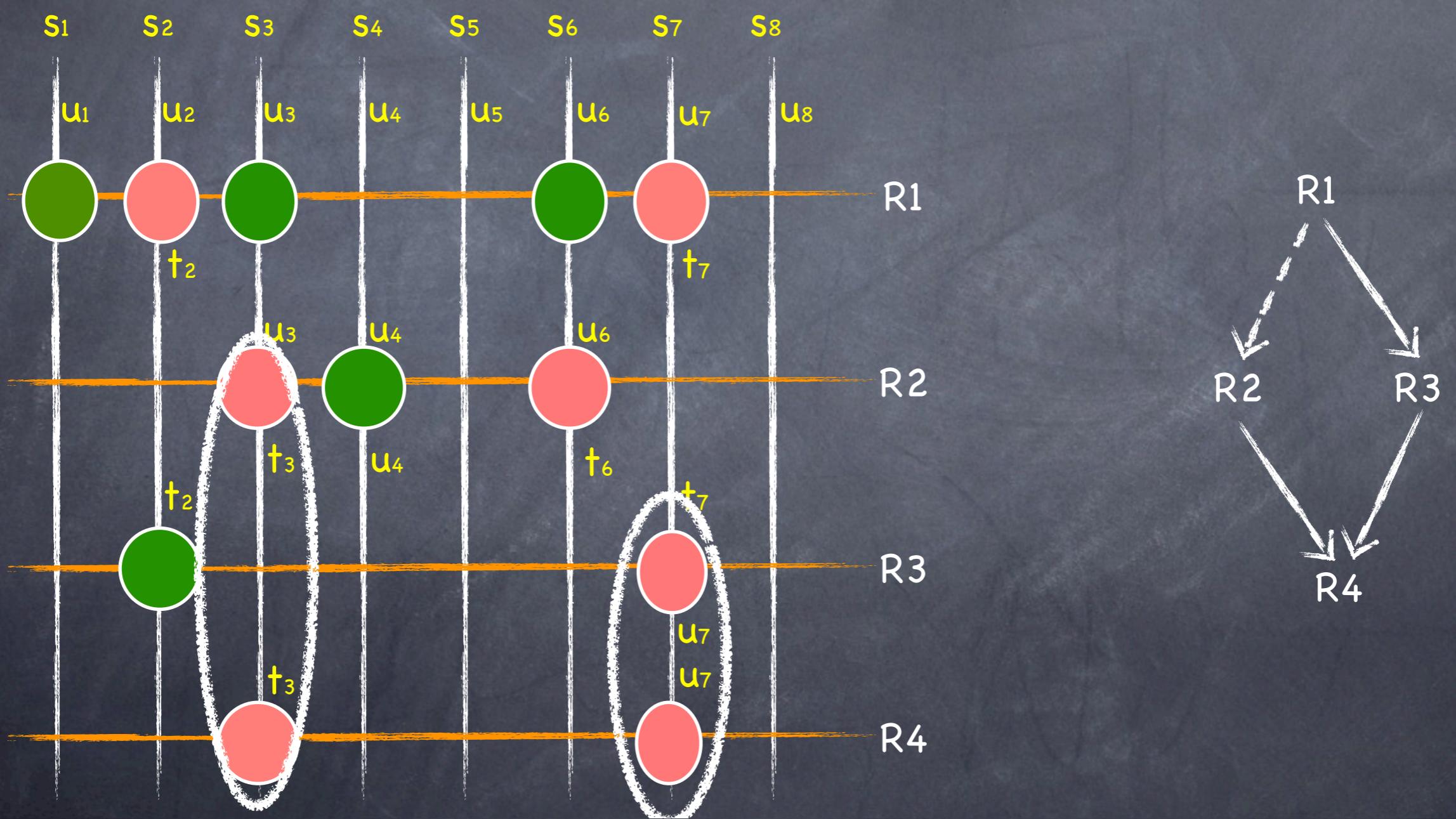
Simple causality analysis



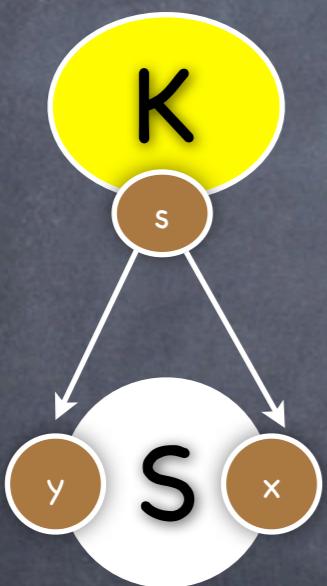
Simple causality analysis



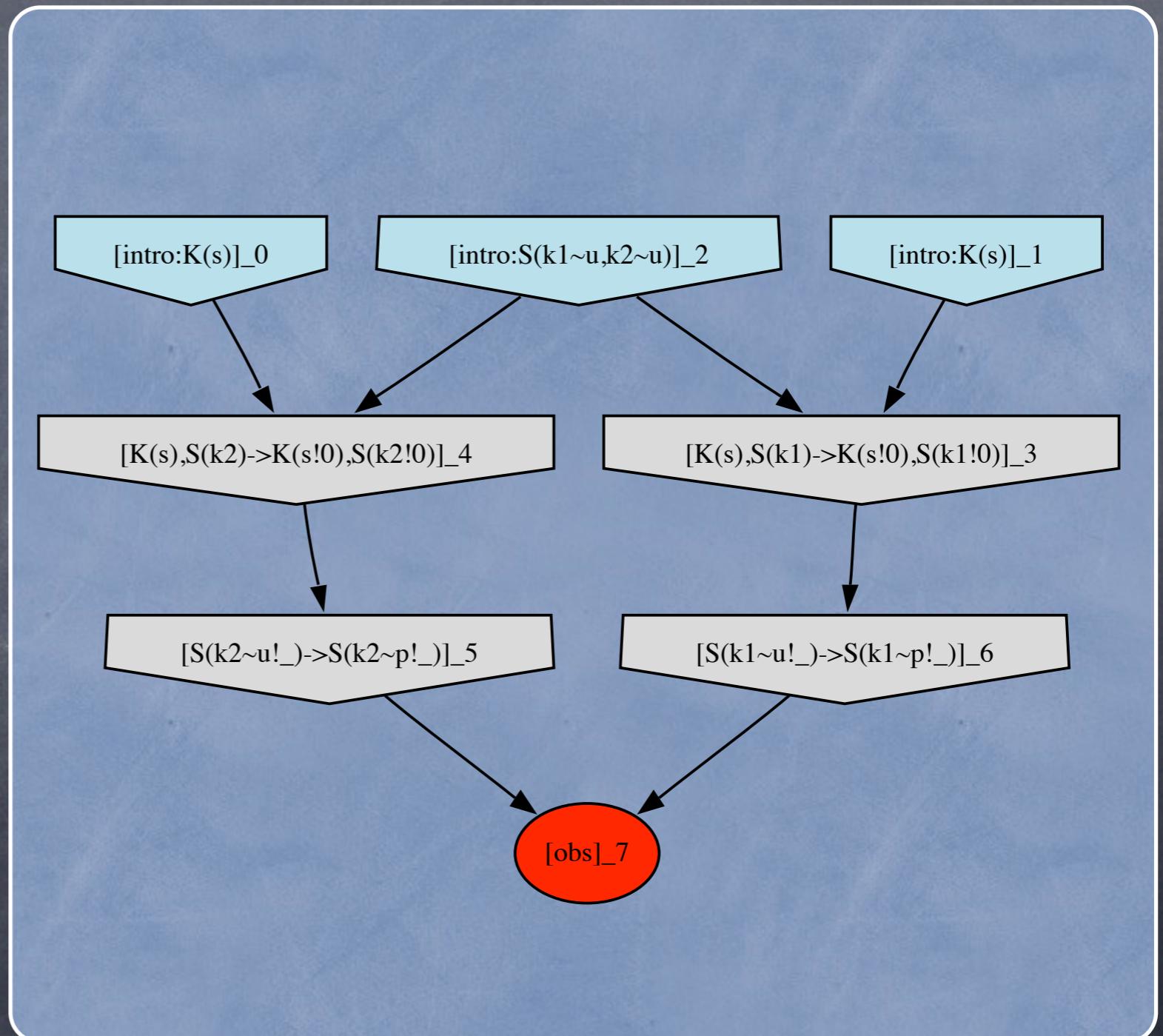
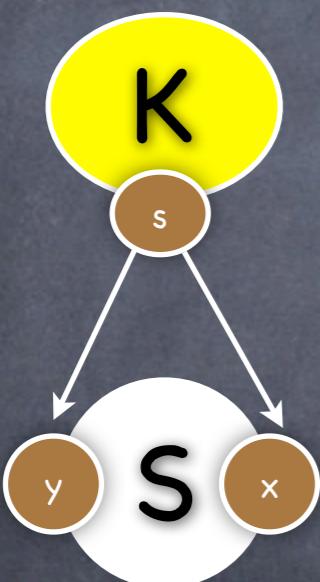
Simple causality analysis



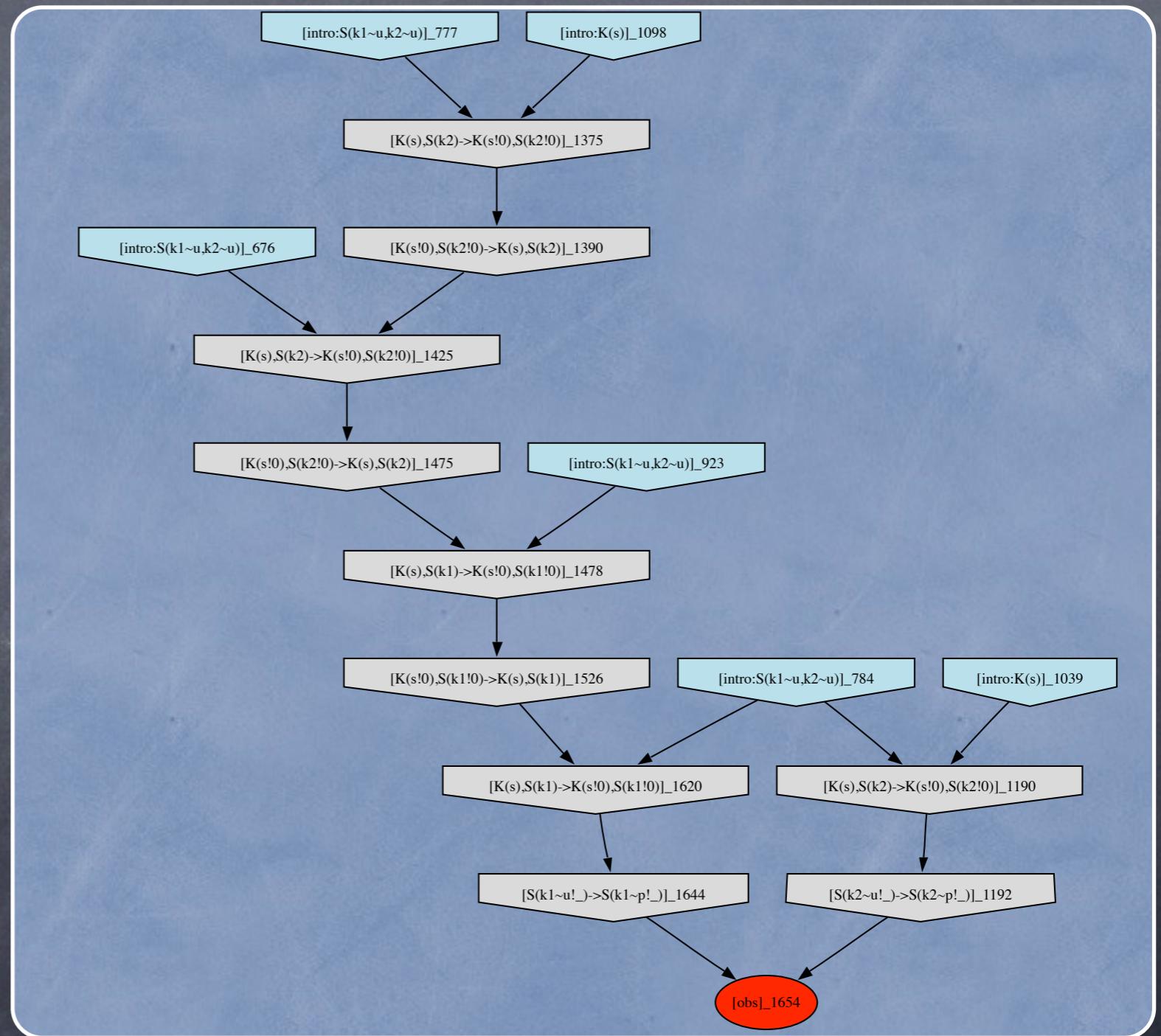
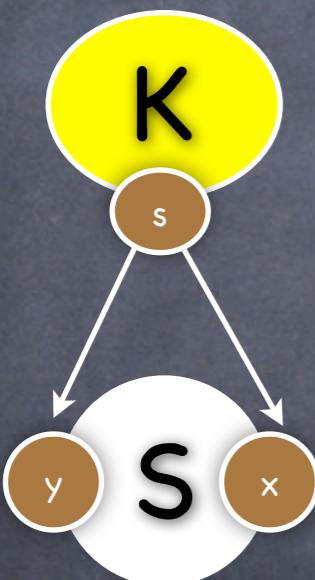
But this is not
satisfactory...



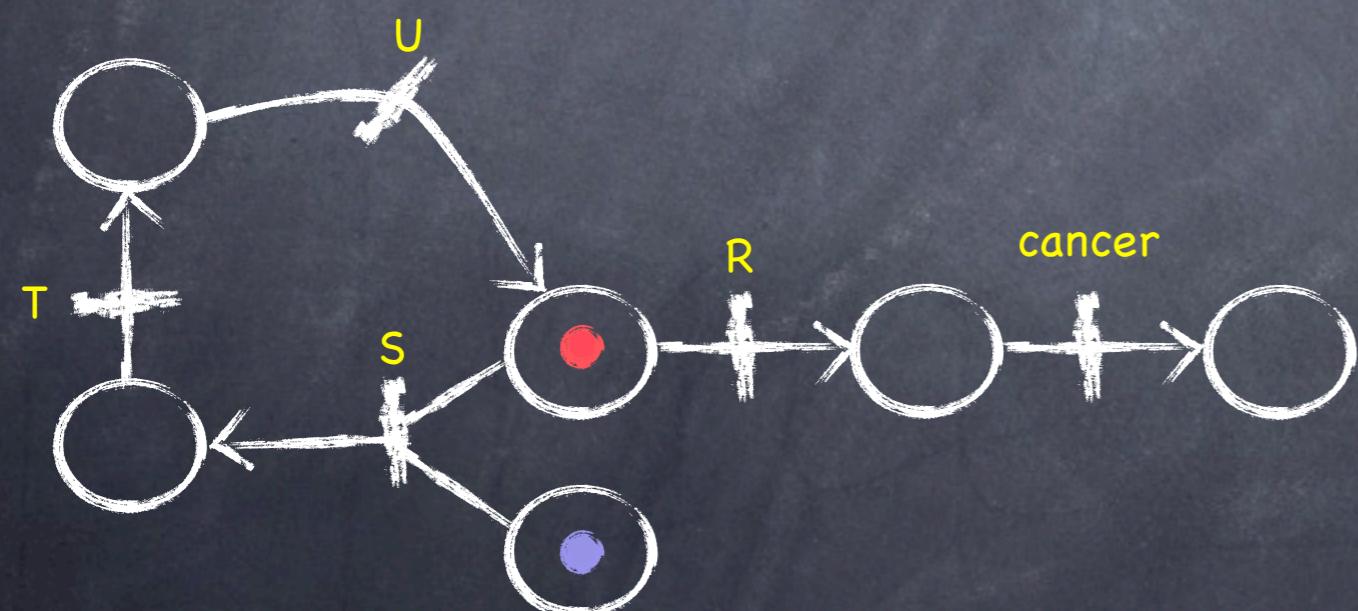
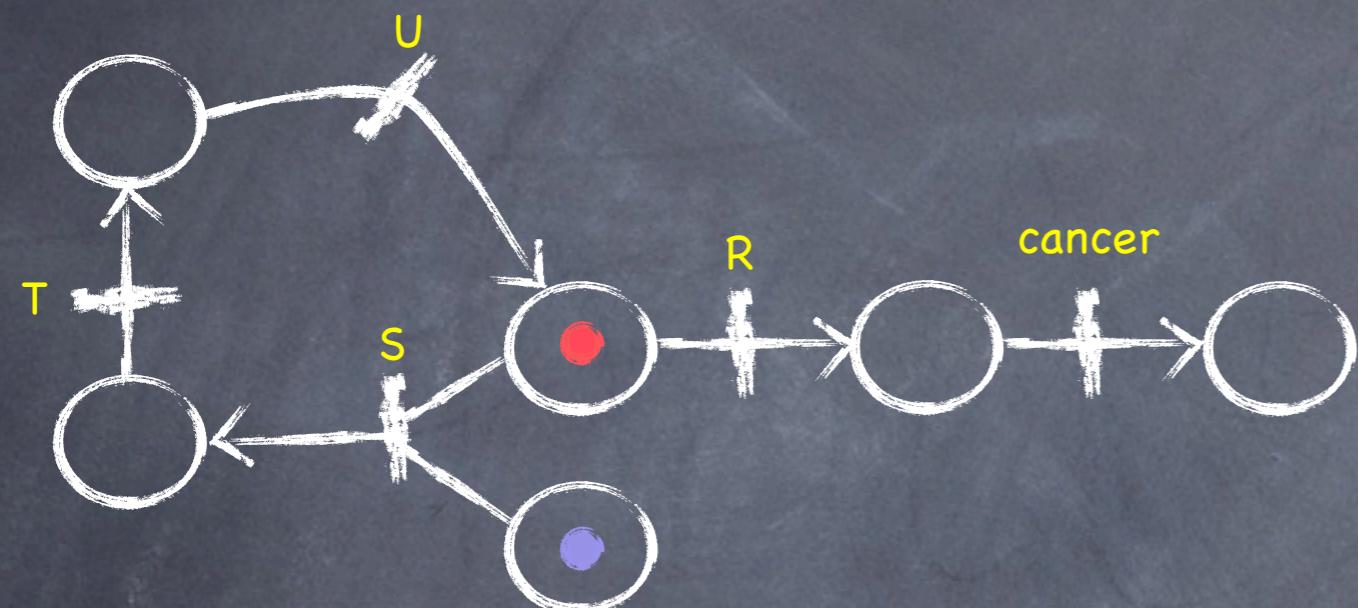
But this is not satisfactory...



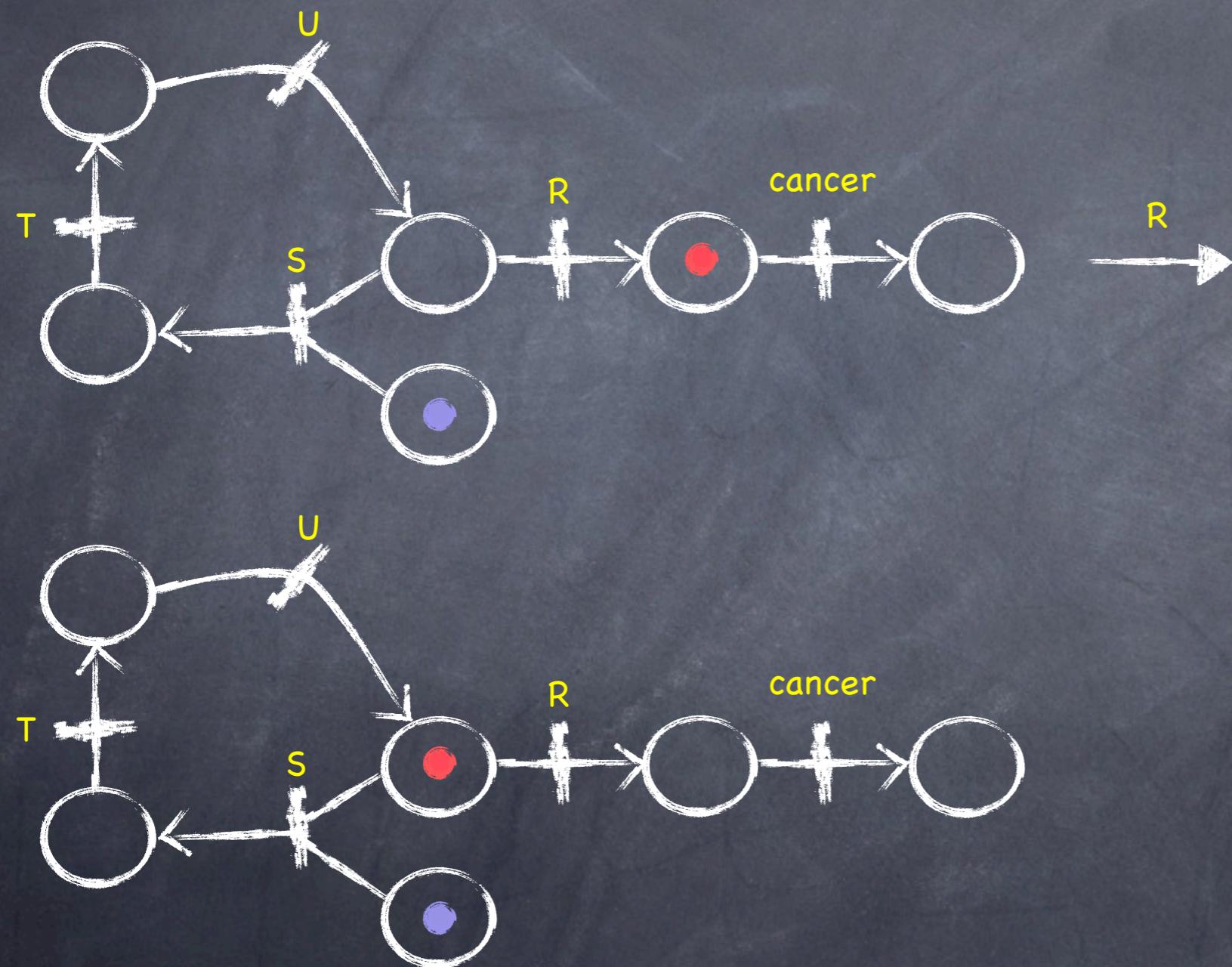
But this is not satisfactory...



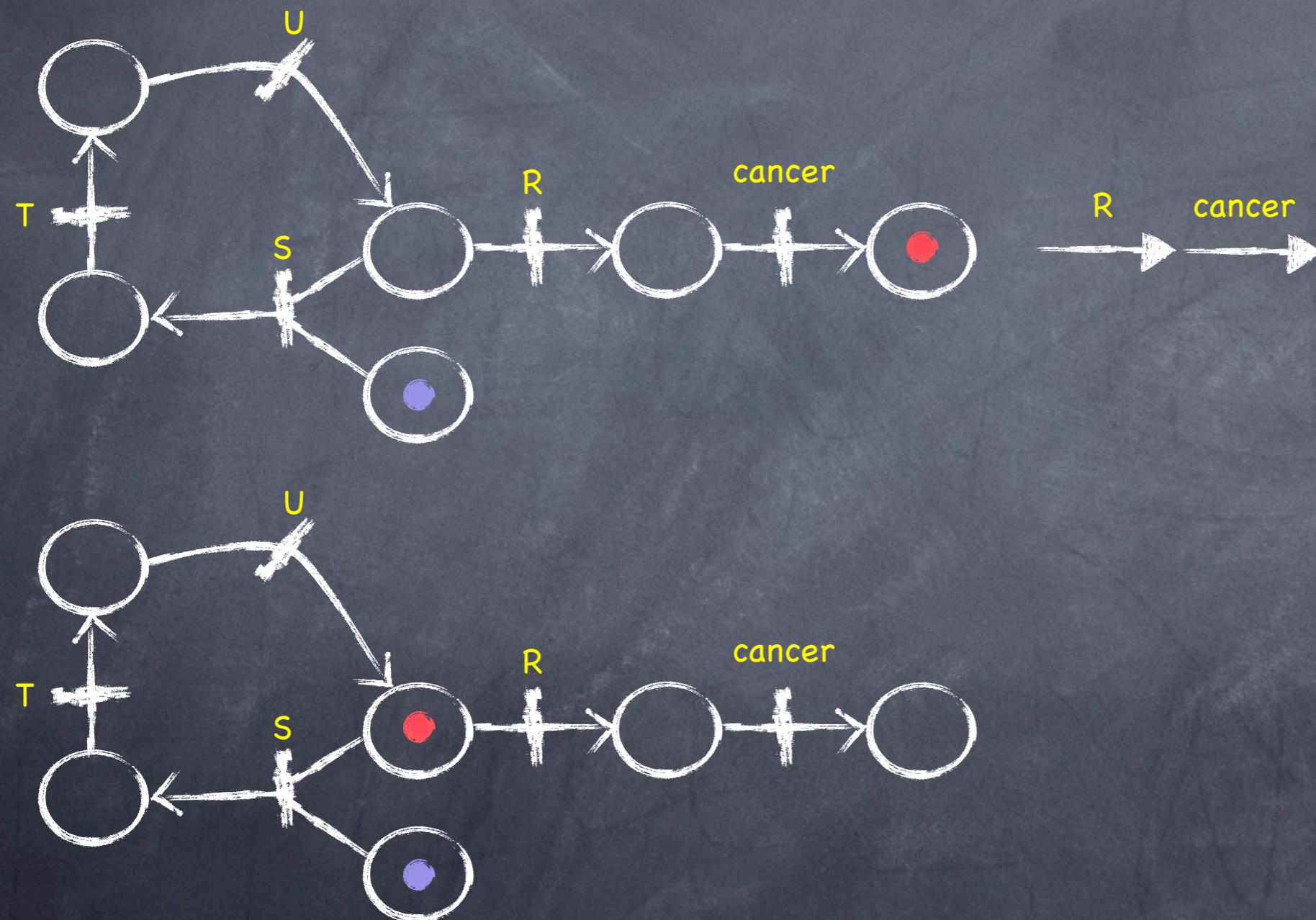
If life was a PT net



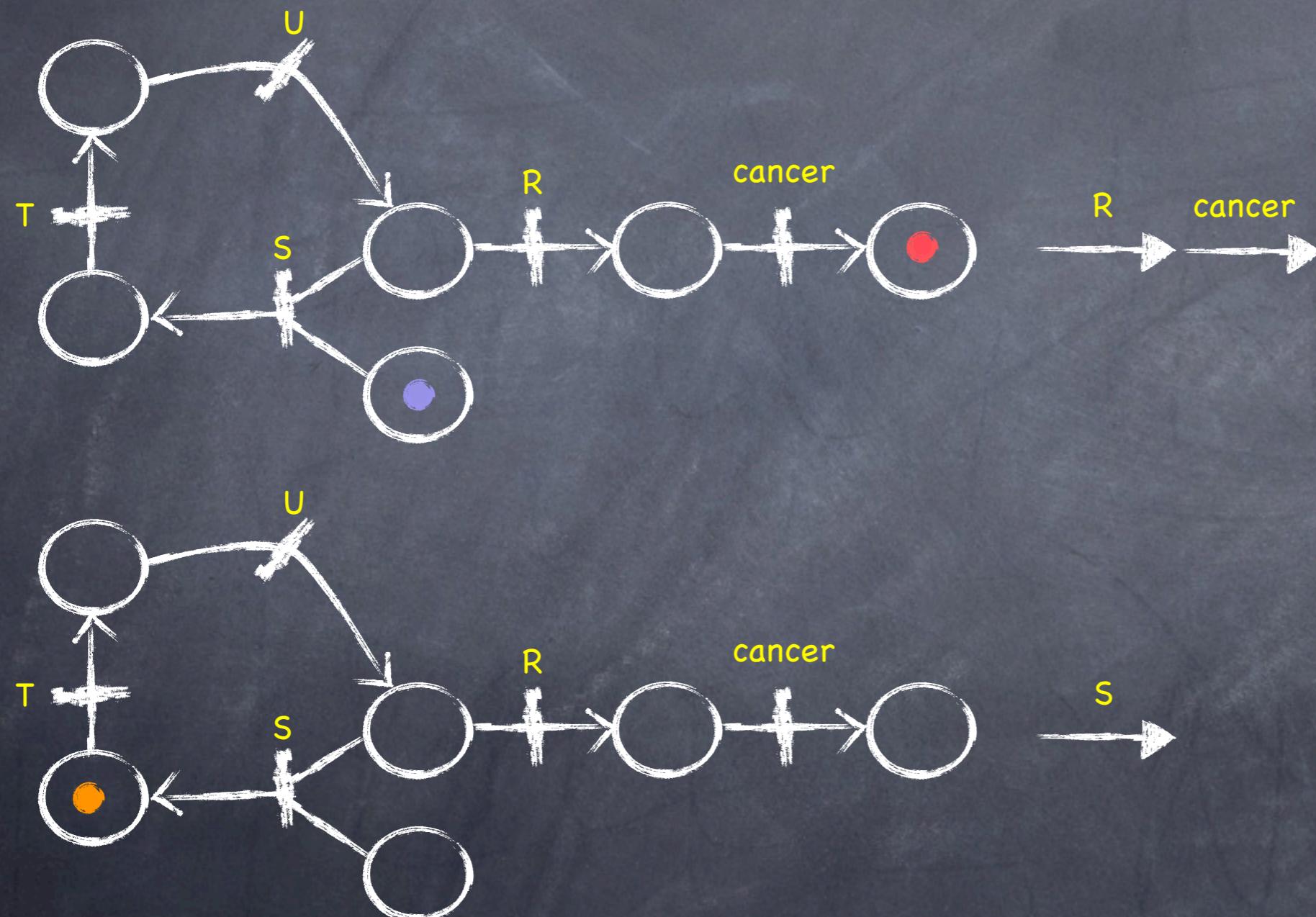
If life was a PT net



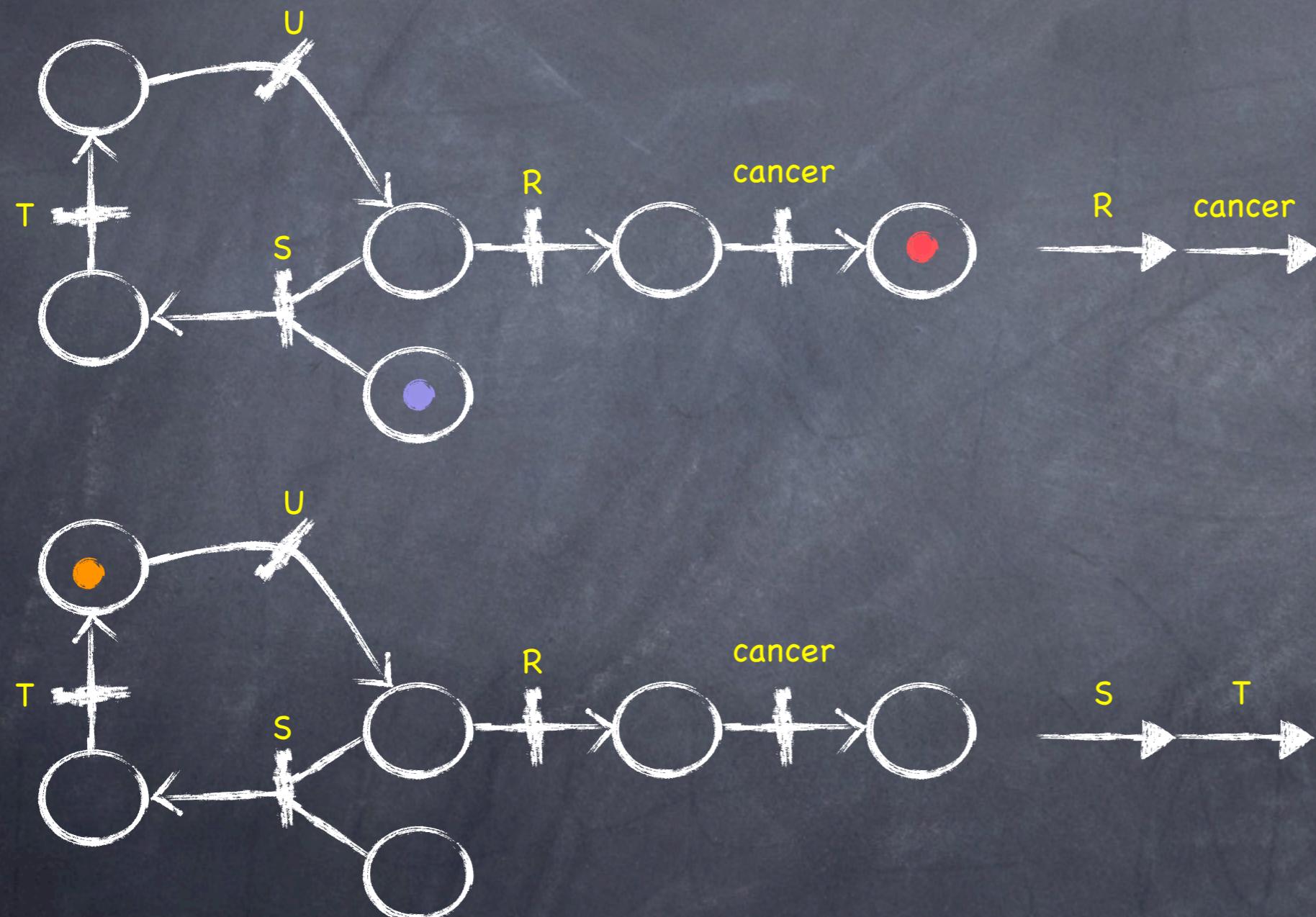
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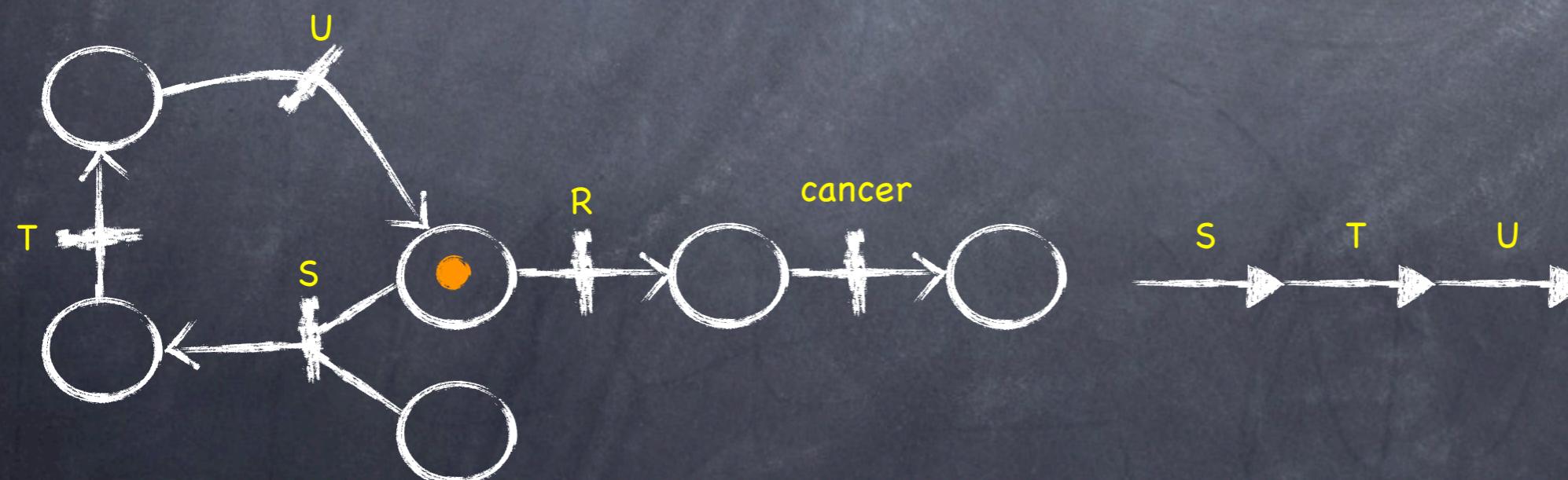
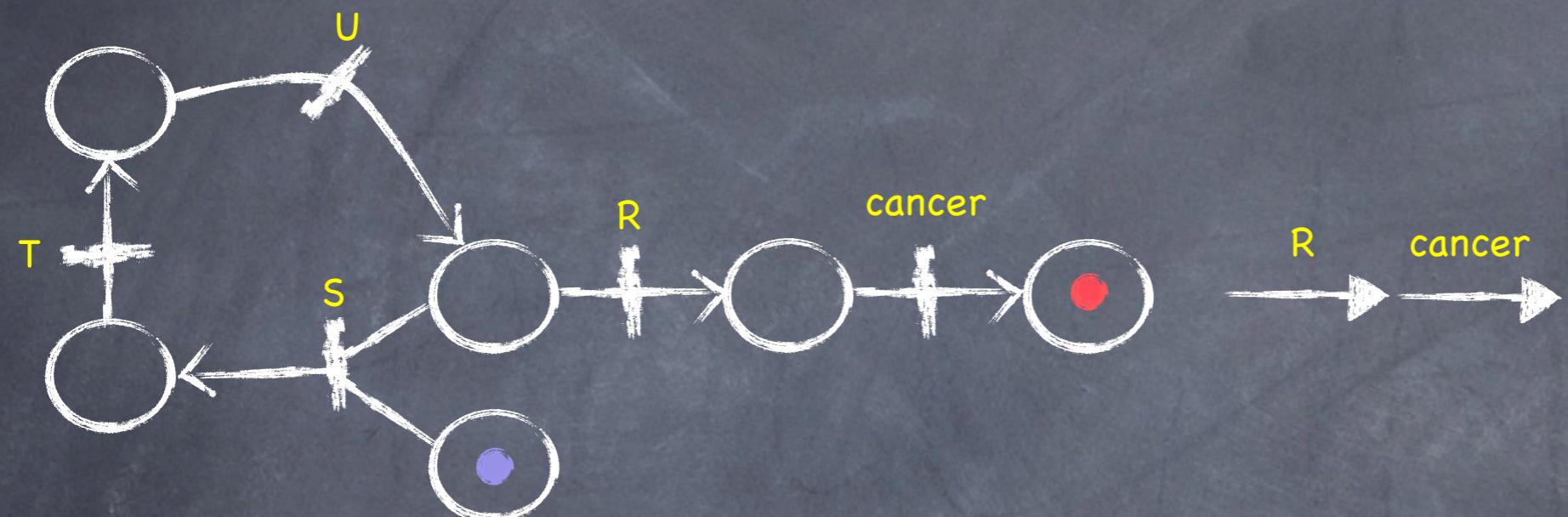
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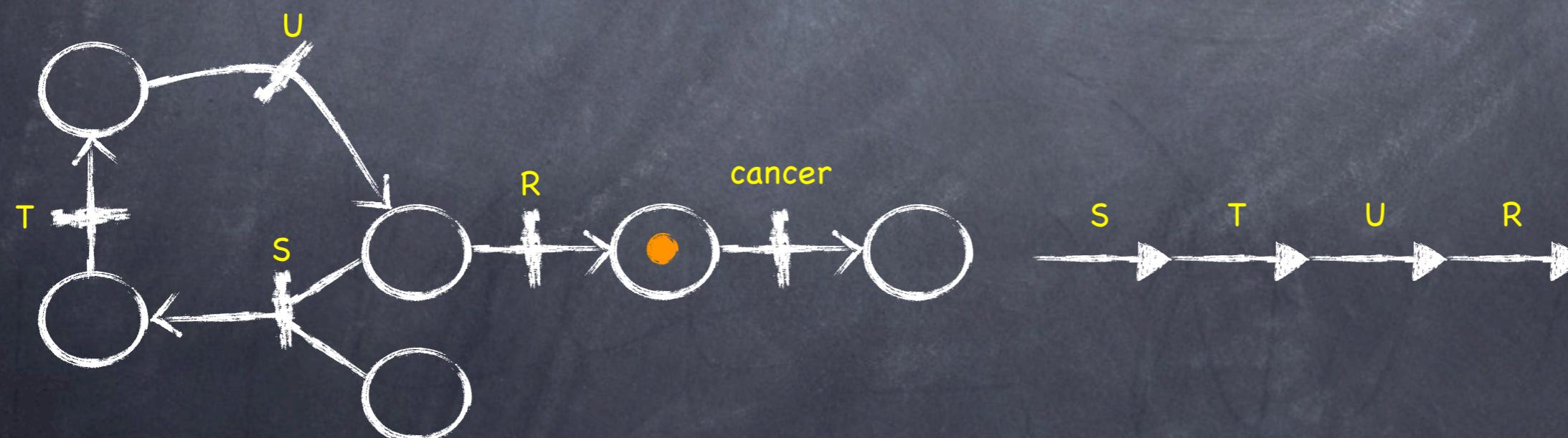
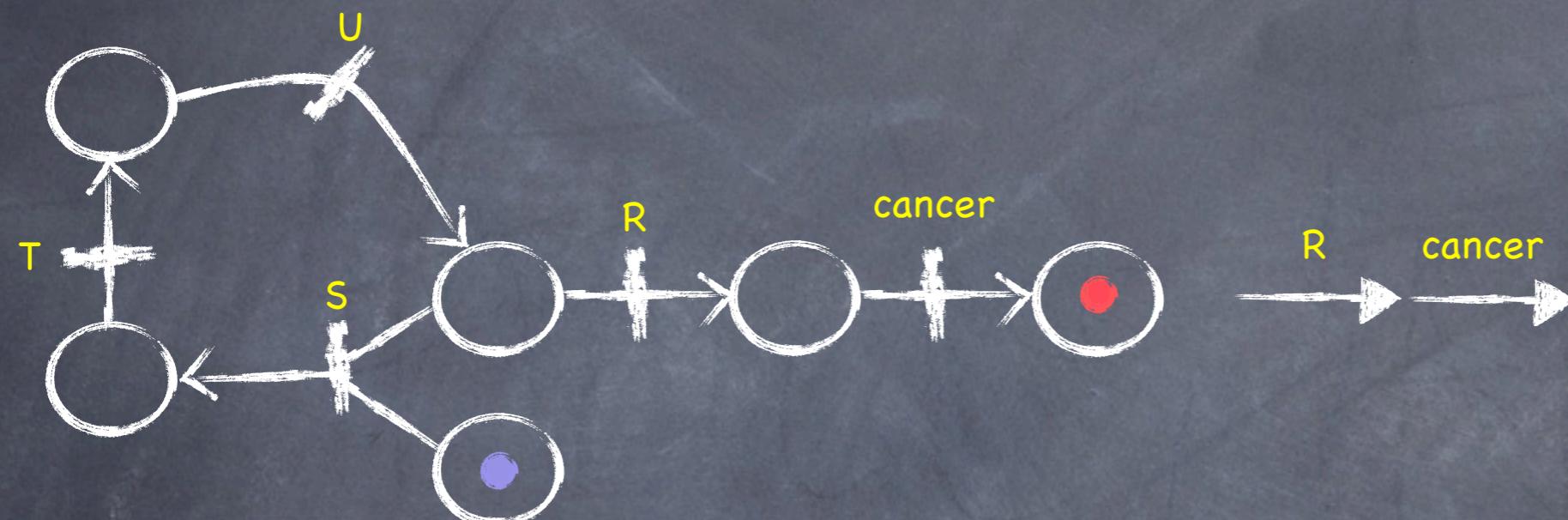
If life was a PT net



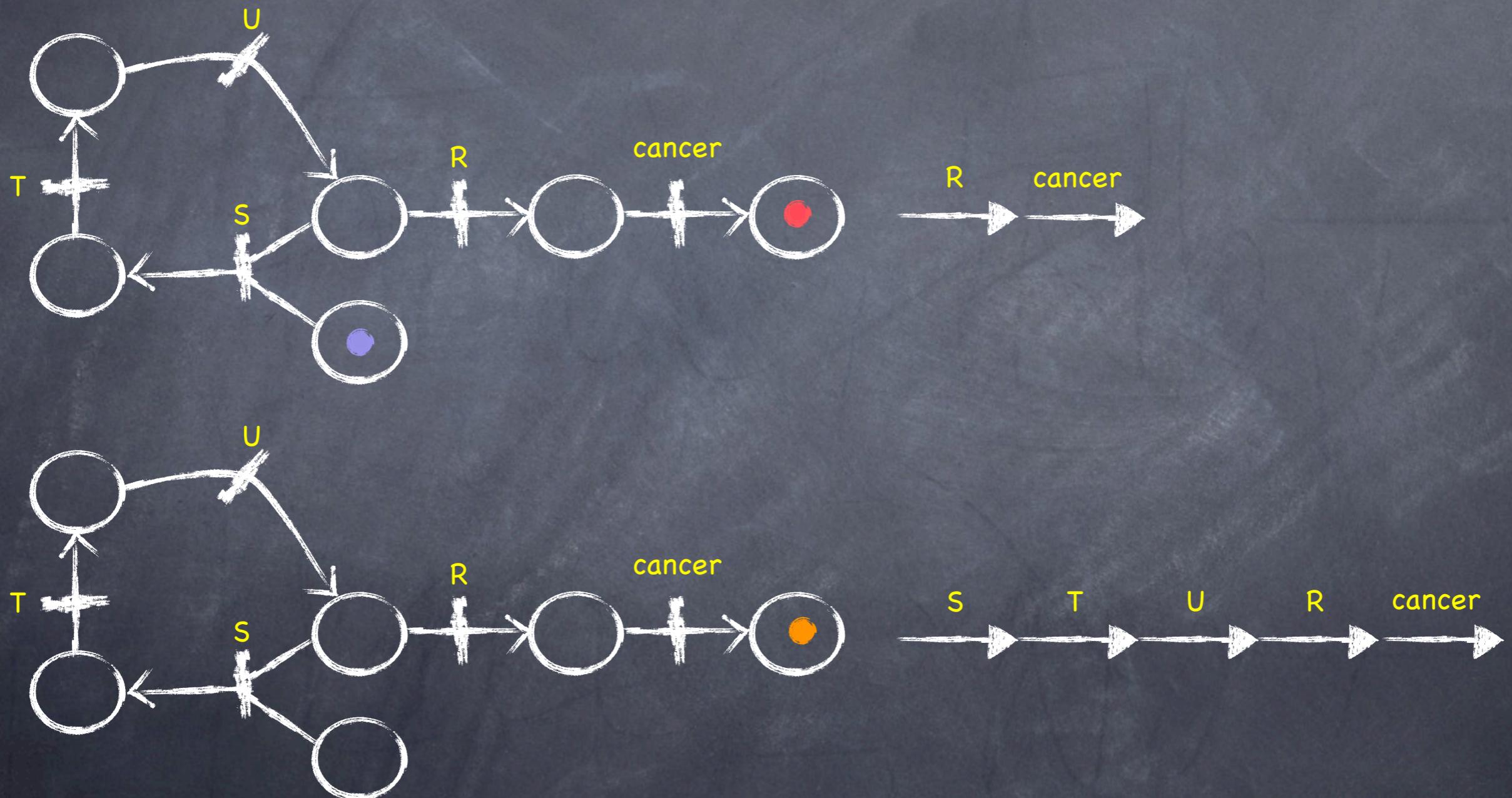
If life was a PT net



If life was a PT net

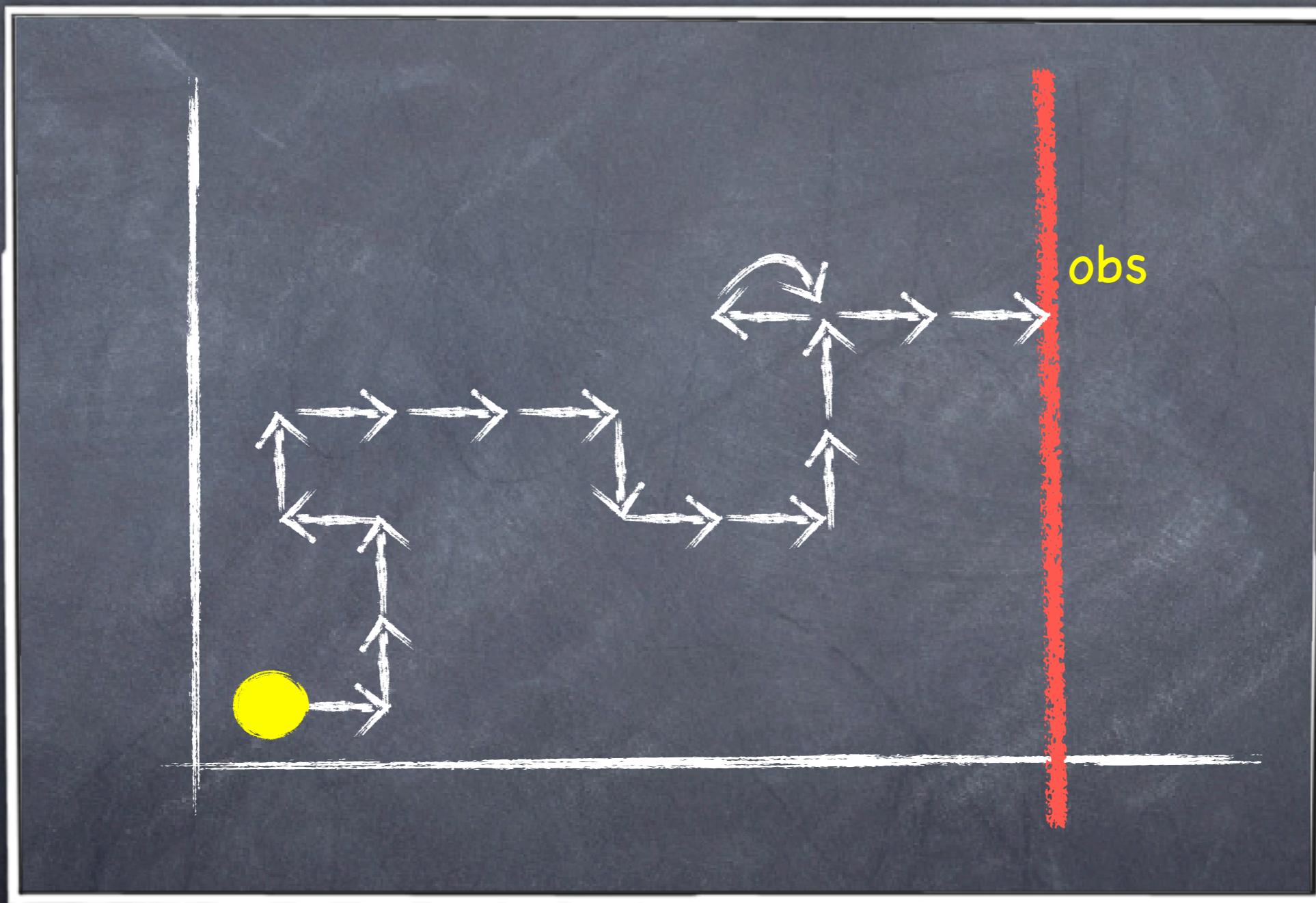


If life was a PT net



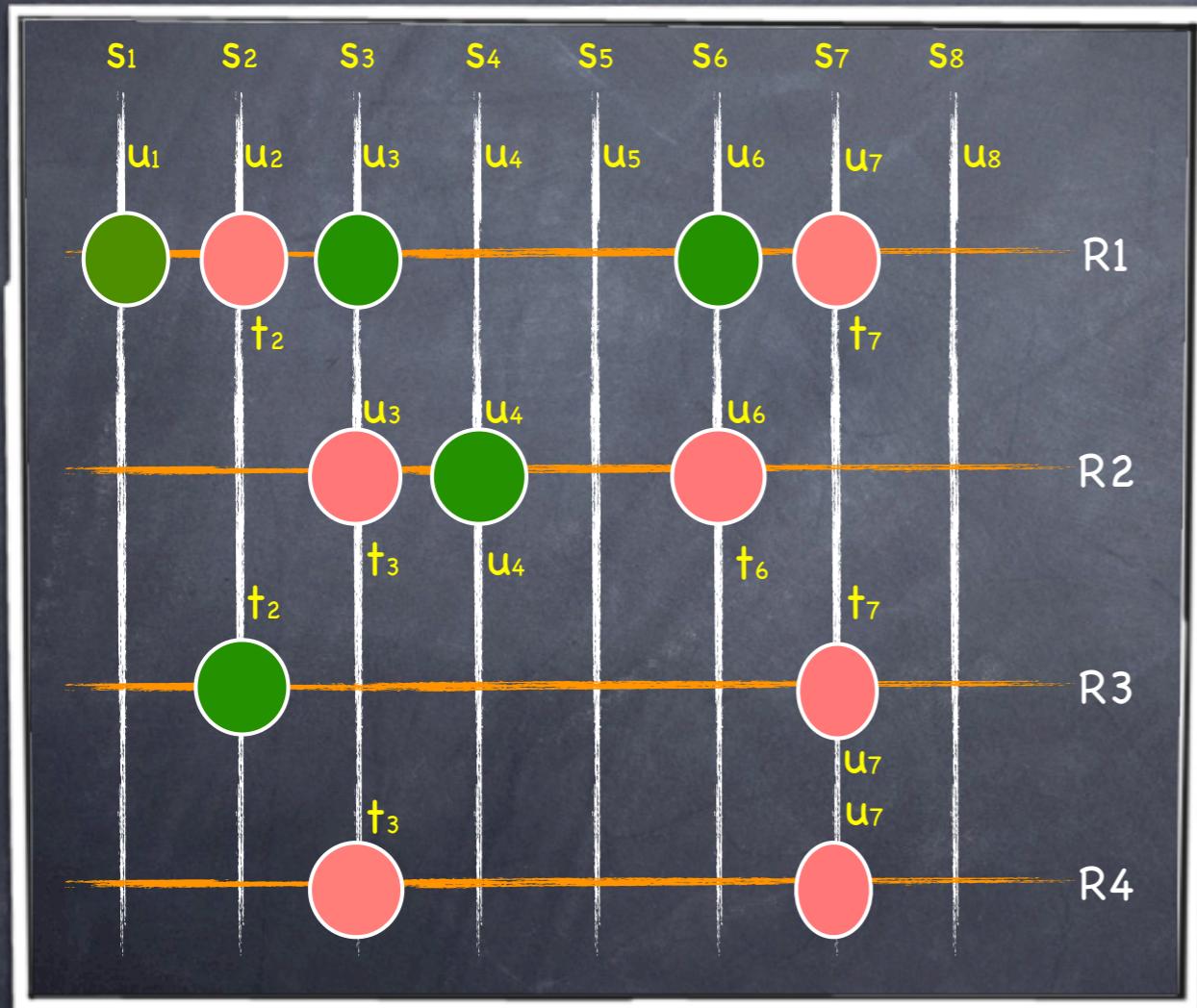
Final states don't match but the difference is not observed

Restricting to what can be observed...



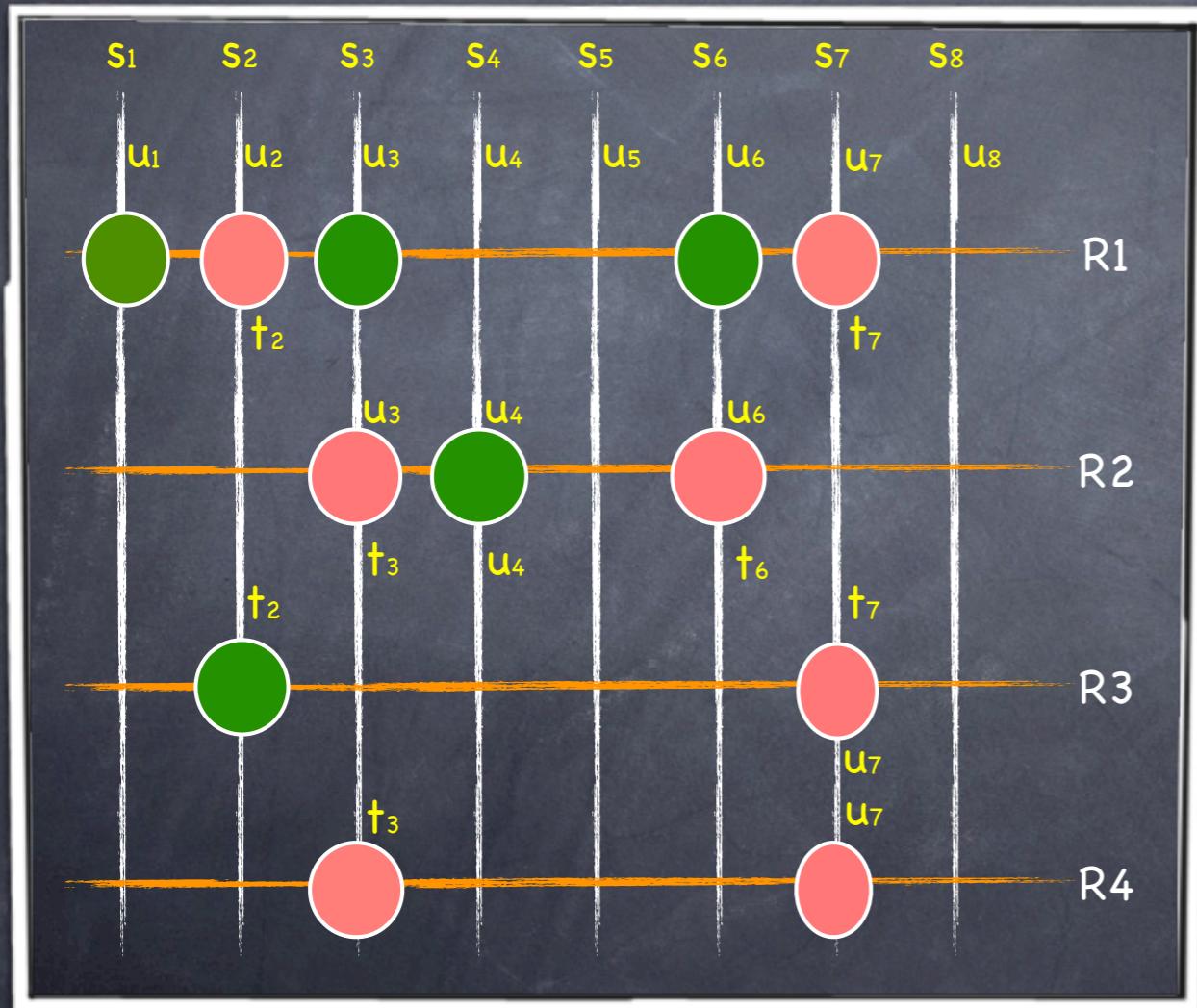
The “knock-out”
property

Is this a causal trace ?



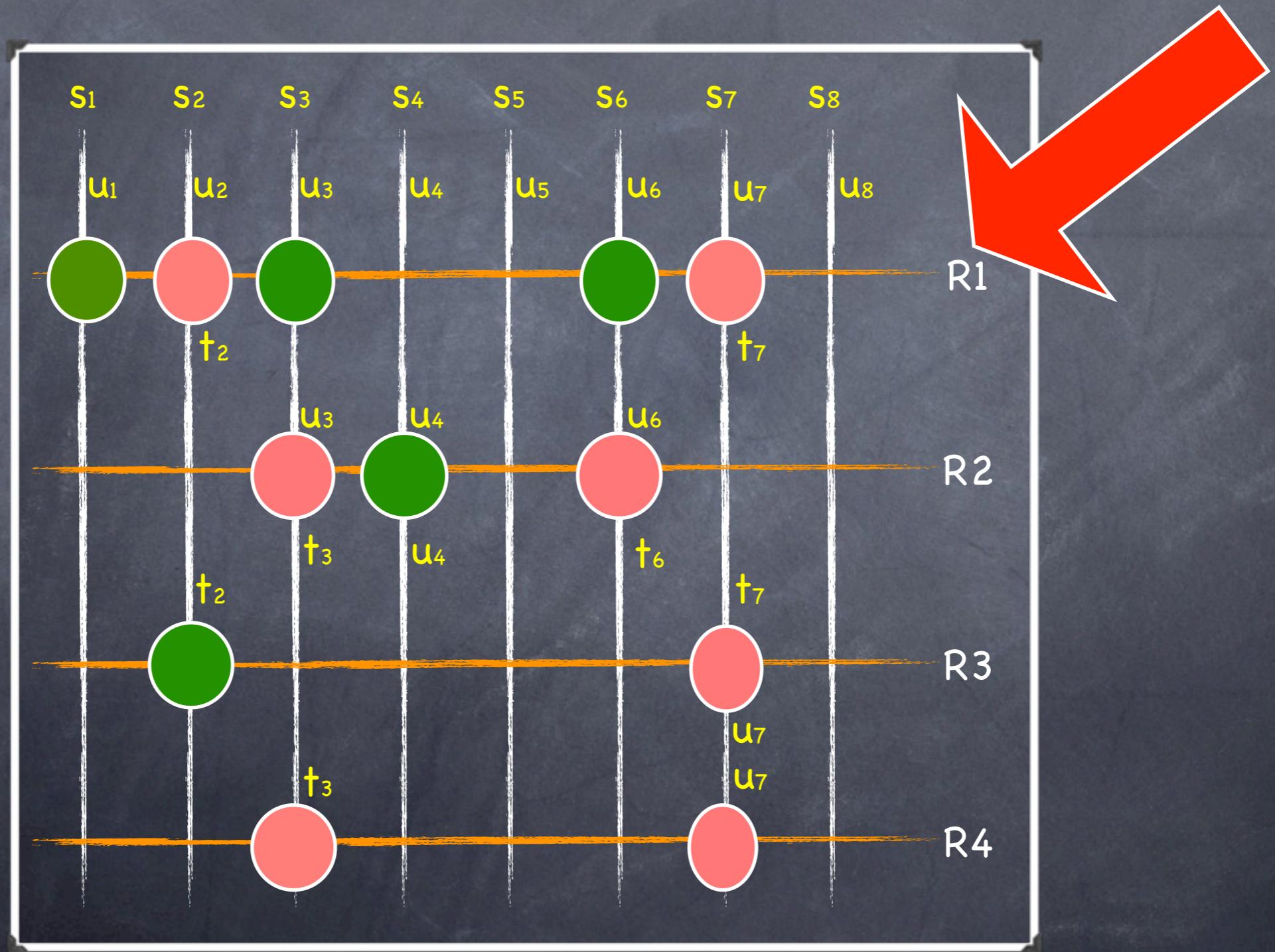
“Yes, because each event is a local cause of the next one”

Is this a causal trace ?

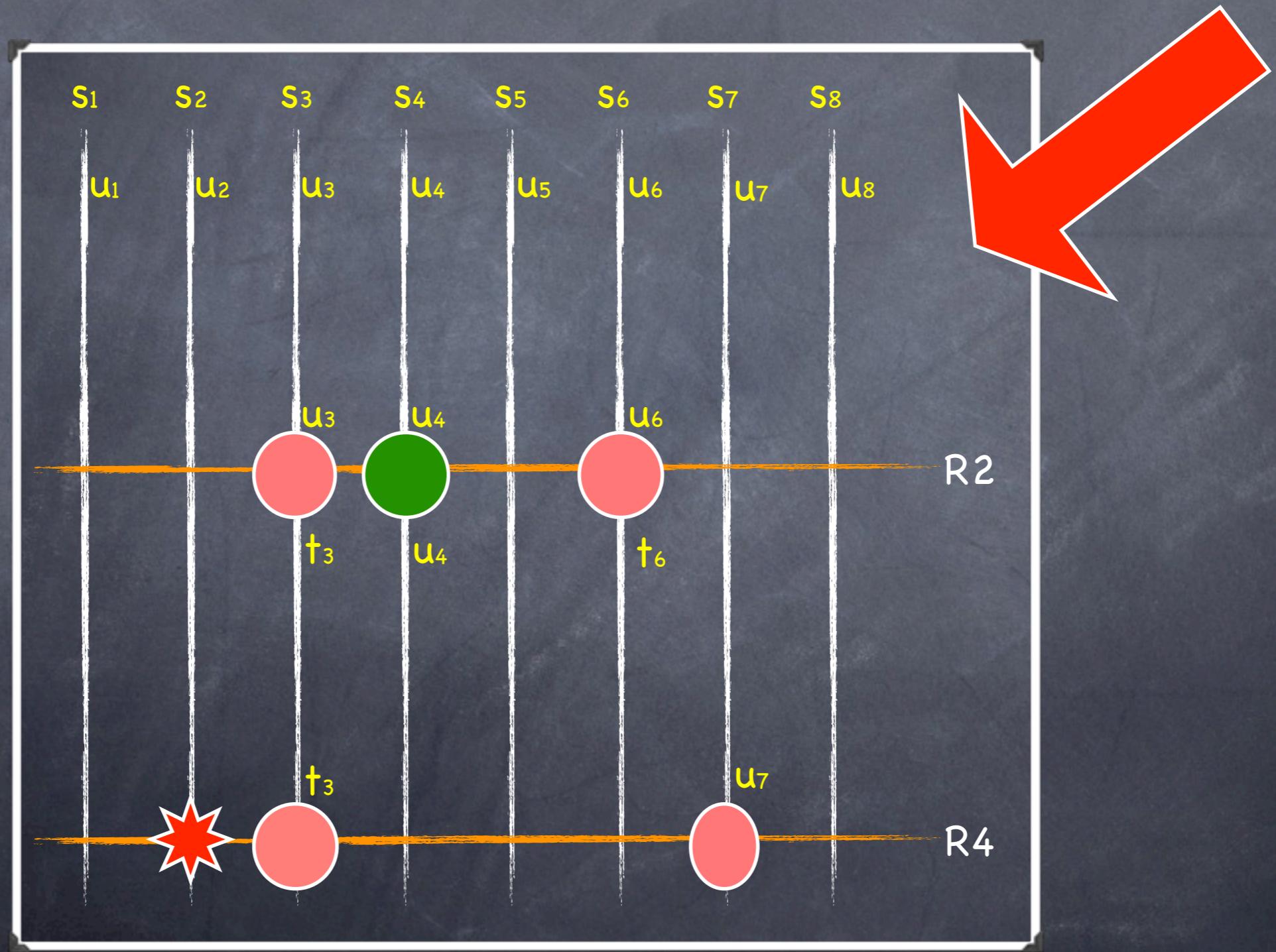


“No, because I can knock-out R1 and still observe R4!”

Knock-out and stabilization



Knock-out and stabilization



Causal Compression

Rule application

A **rule** denote an action:

$$S \xrightarrow{a} T$$

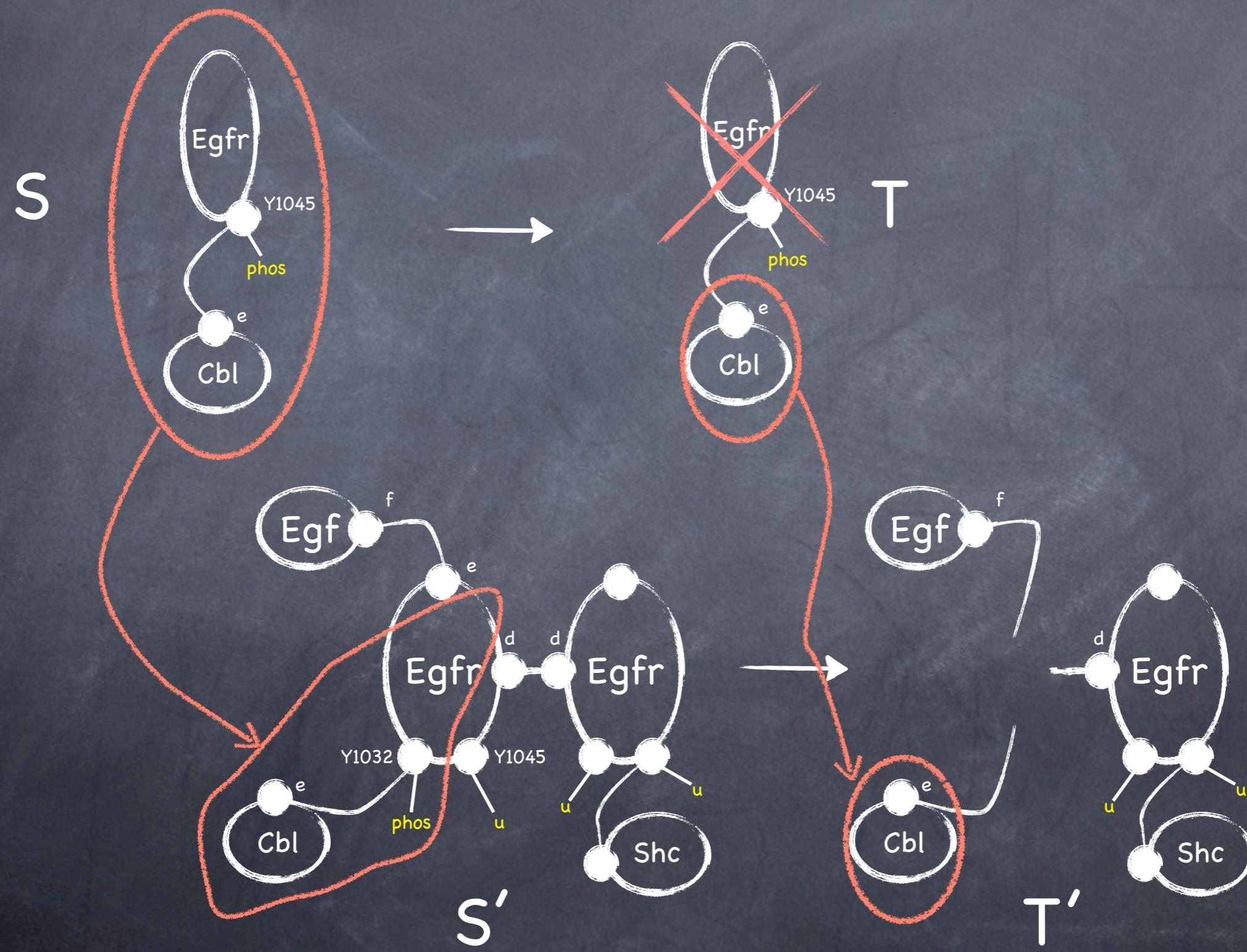
An **application** of such rule is determined by an embedding:

$$S \xrightarrow{e} S'$$

Its **effect** is given by the pushout:

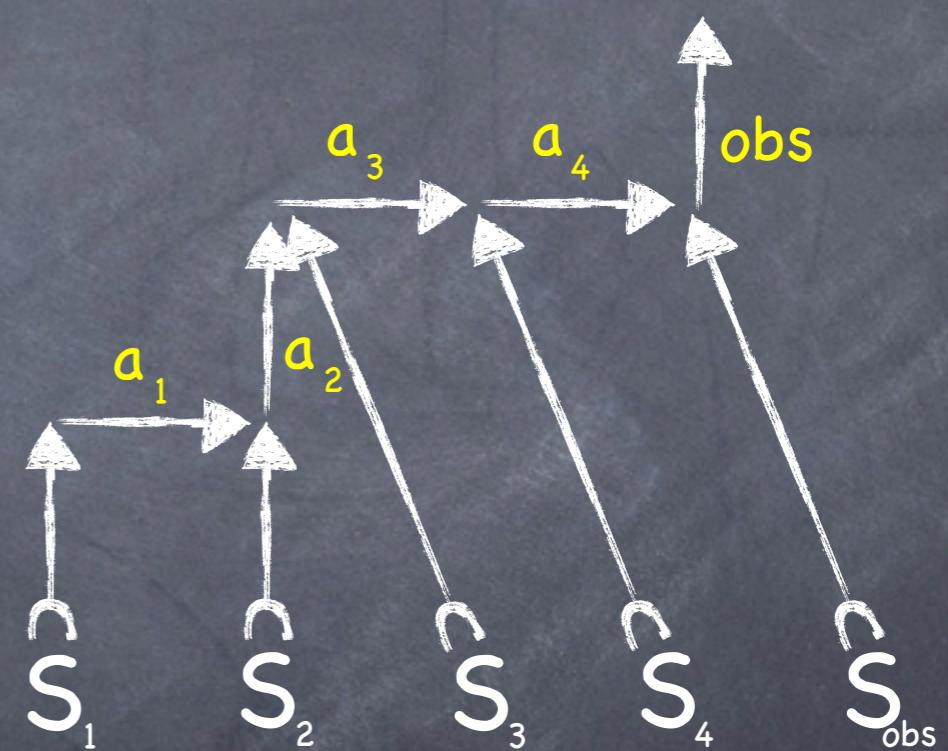
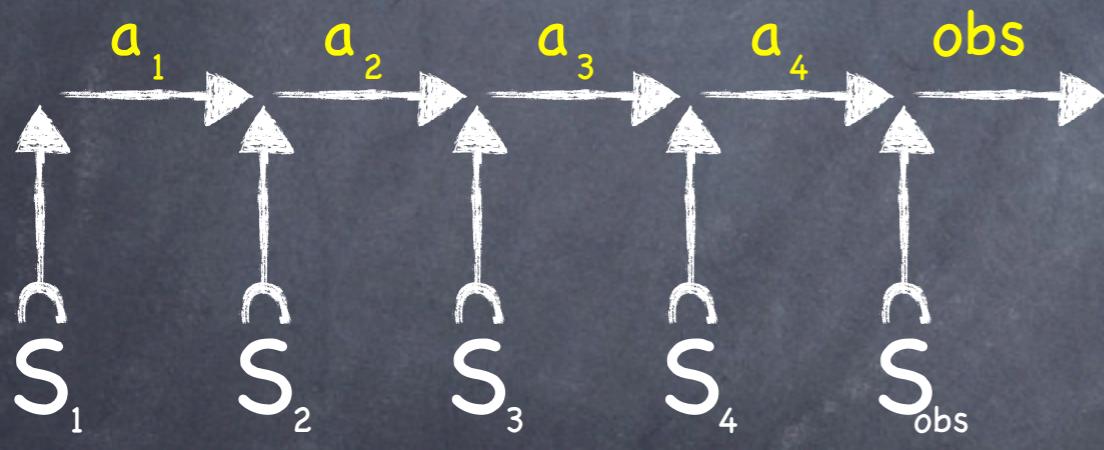
$$\begin{array}{ccc} S & \xrightarrow{a} & T \\ e \downarrow & & \downarrow e' \\ S' & \xrightarrow{a'} & T \end{array}$$

For instance



Simple compression

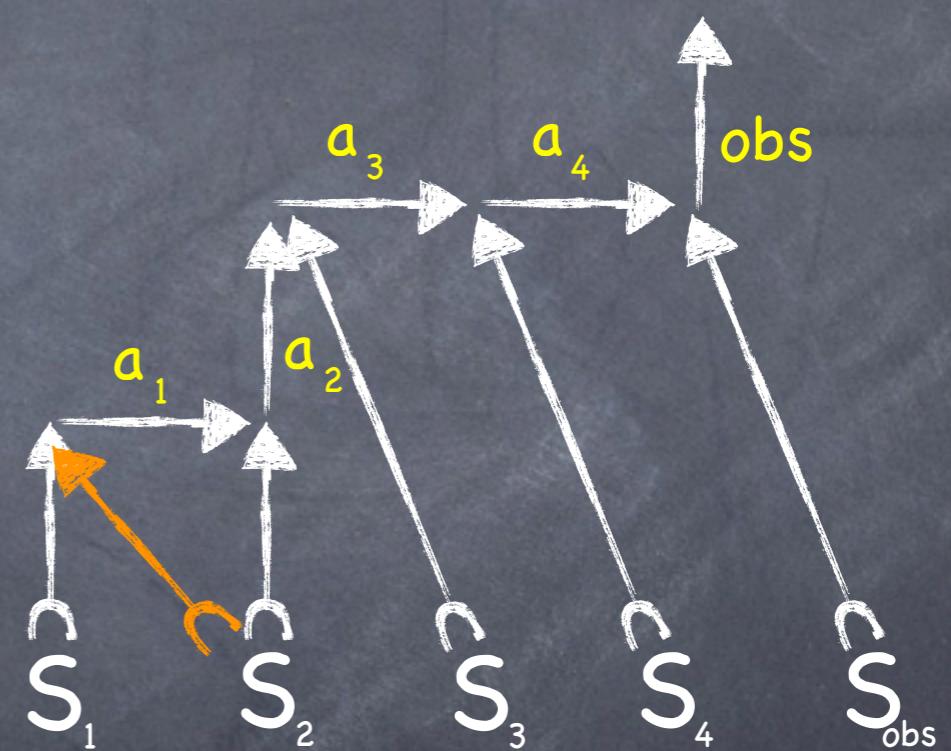
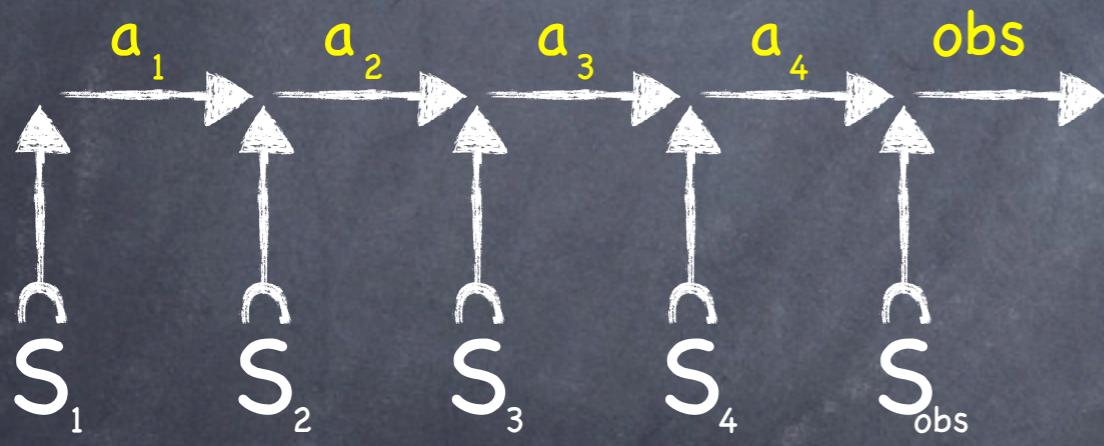
A trace (obtained by simulation)



Assume the above partition of arrows is given by an oracle

Simple compression

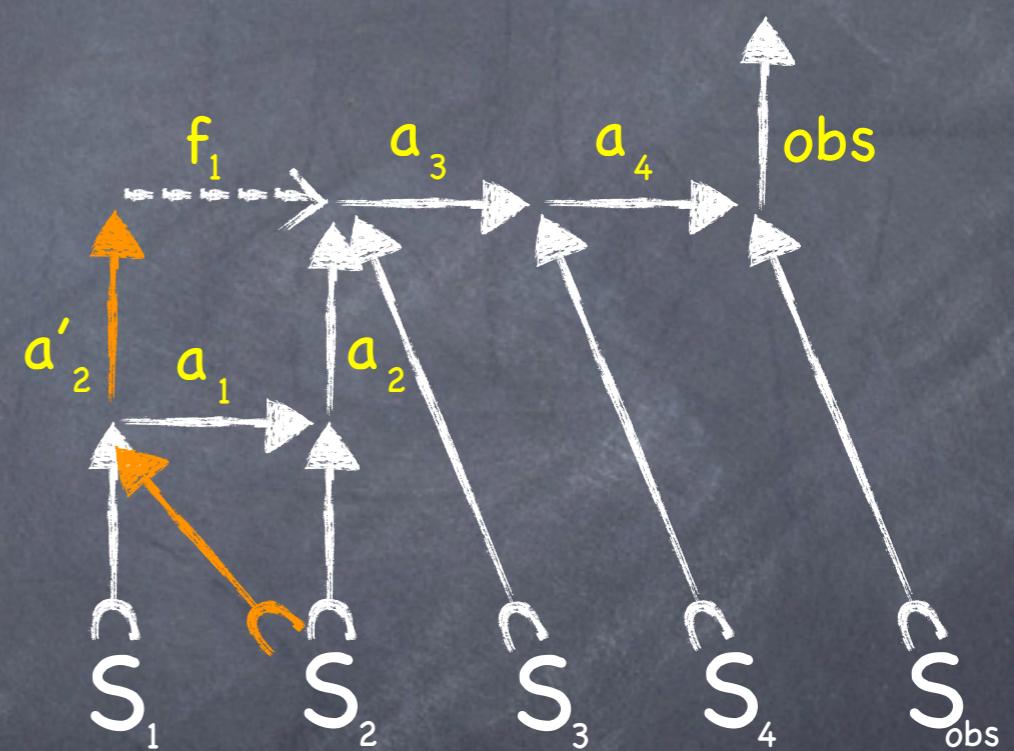
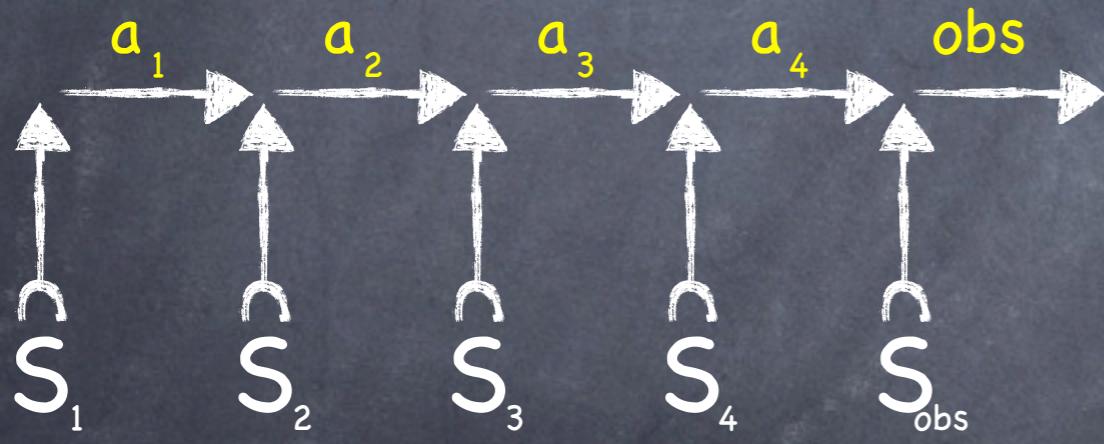
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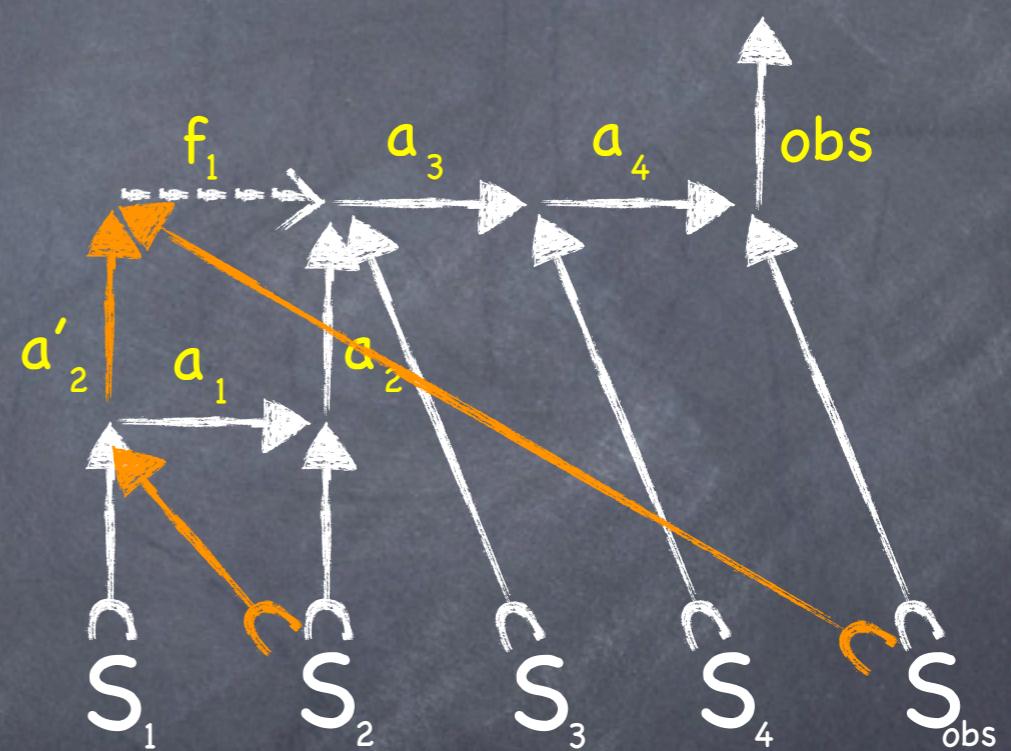
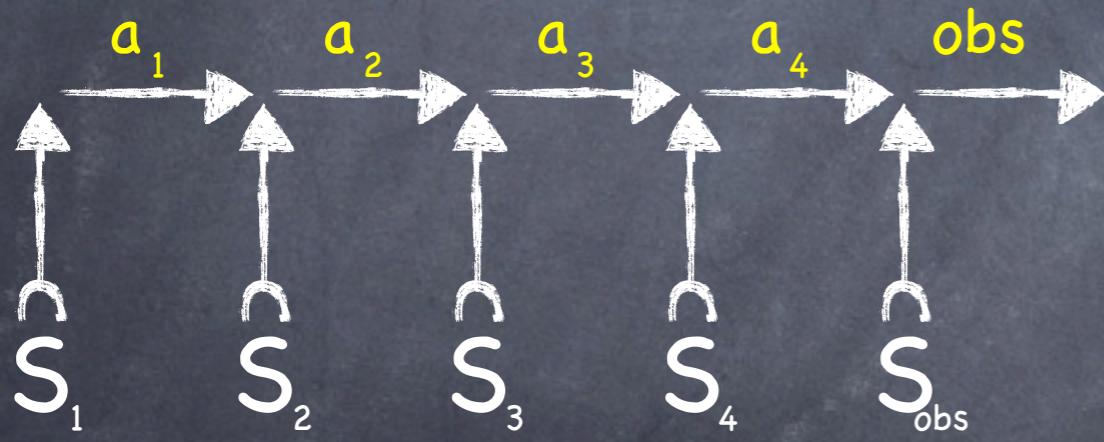
A trace (obtained by simulation)



Assume the above partition of arrows is given by an oracle

Simple compression

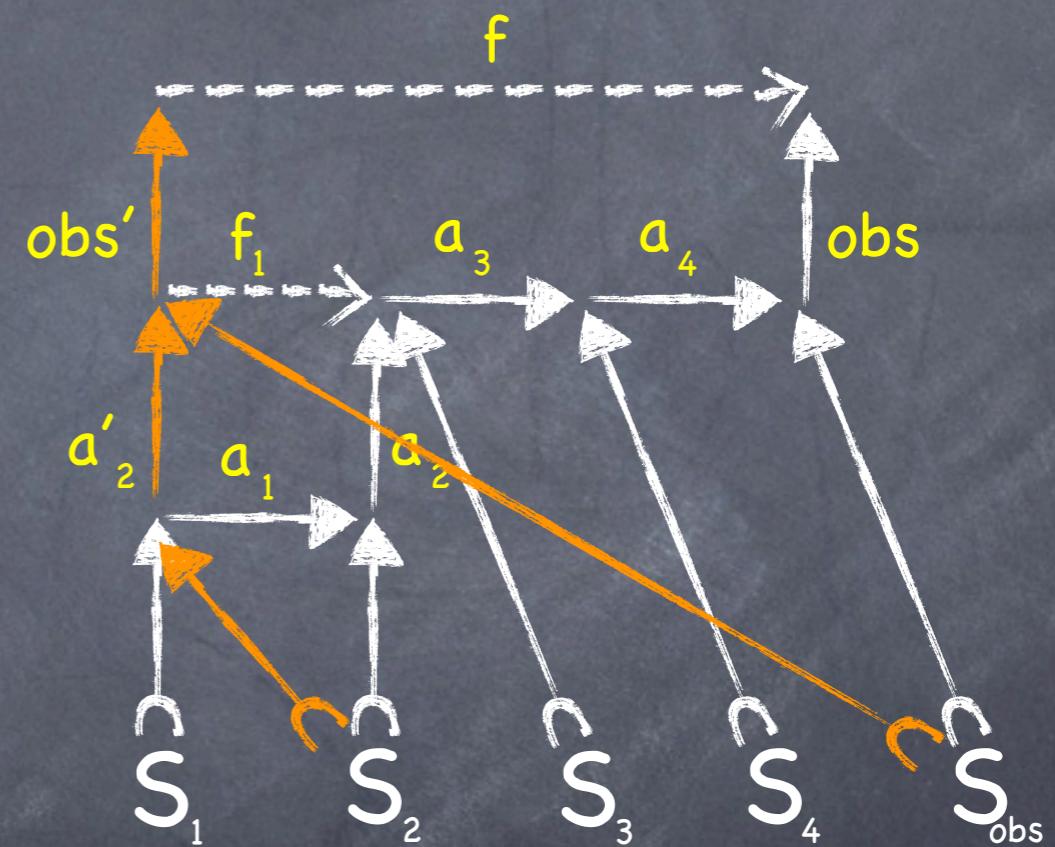
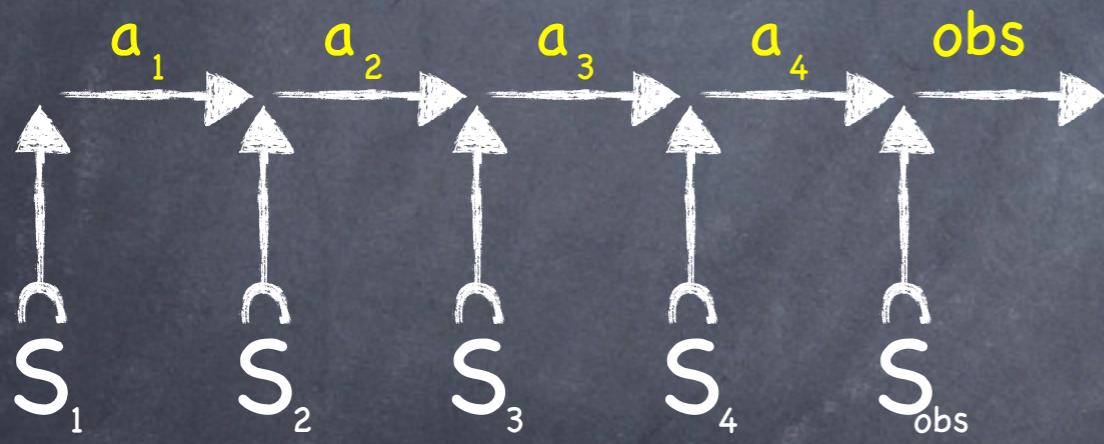
A trace (obtained by simulation)



Assume the above partition of arrows is given by an oracle

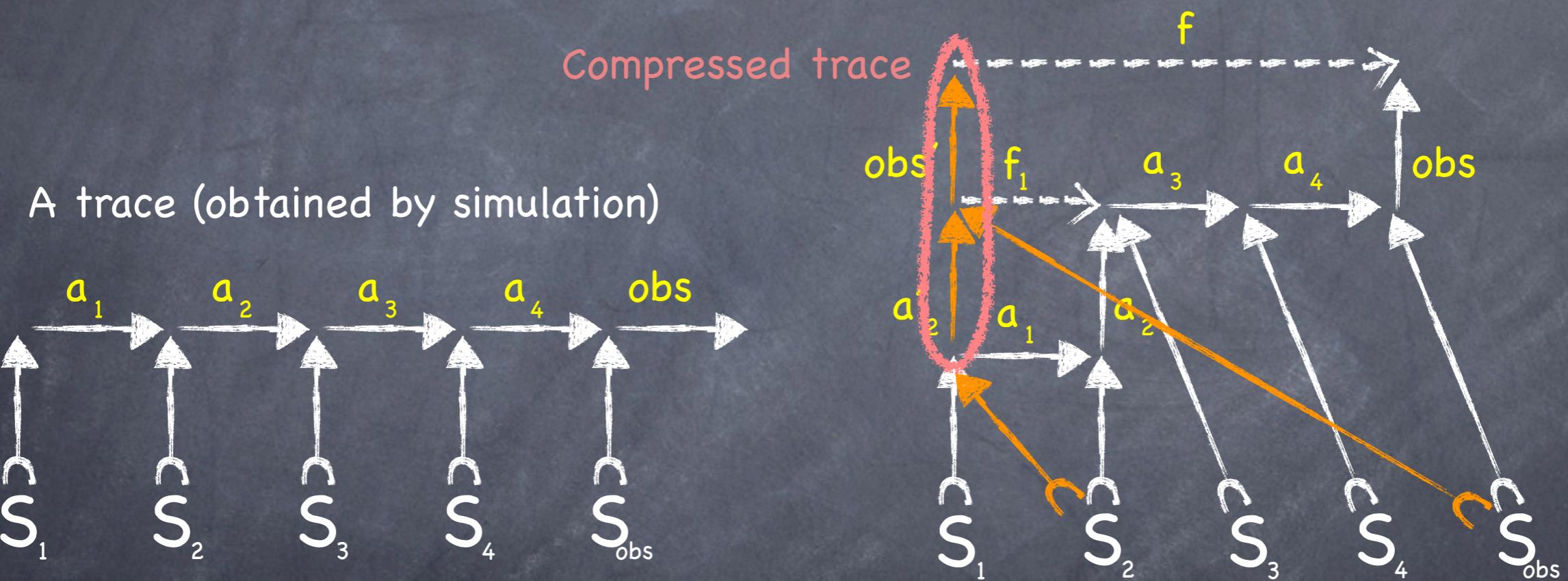
Simple compression

A trace (obtained by simulation)



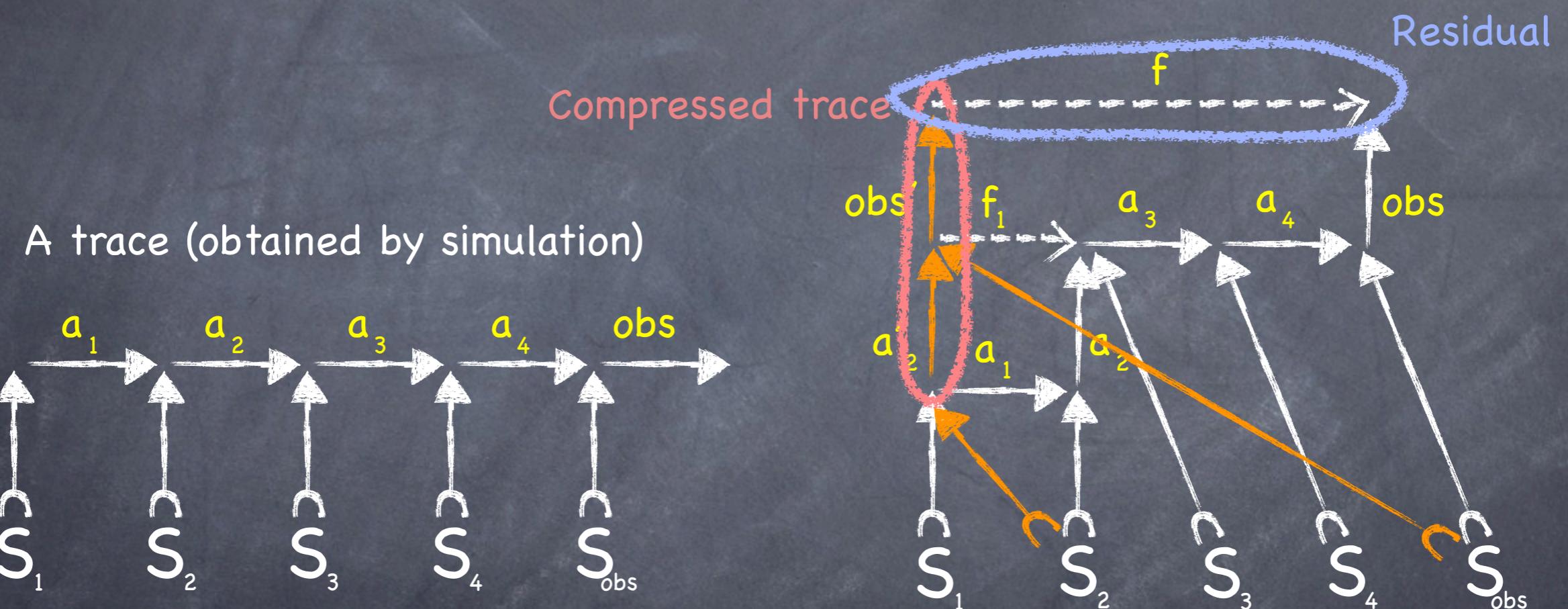
Assume the above partition of arrows is given by an oracle

Simple compression



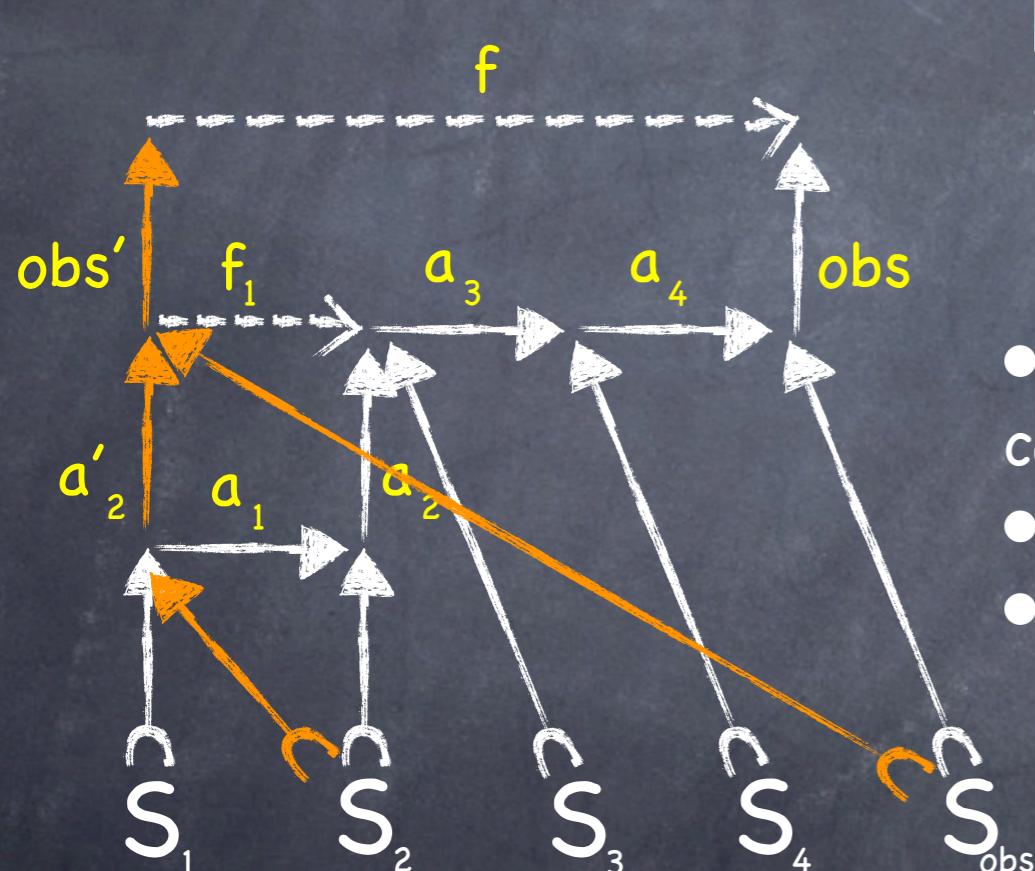
Assume the above partition of arrows is given by an oracle

Simple compression



Assume the above partition of arrows is given by an oracle

Generalized commutation

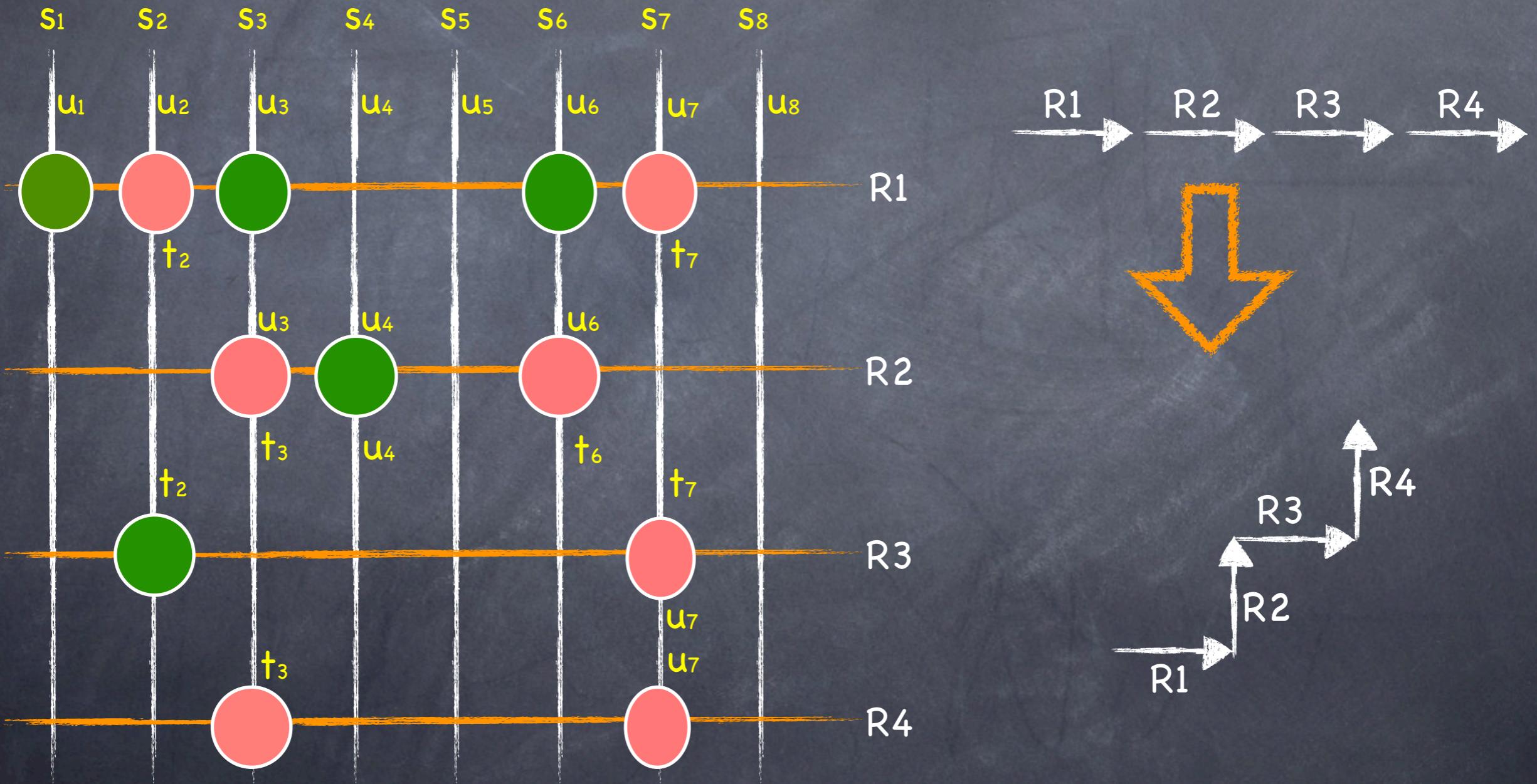


Looking at f gives us some information about compression:

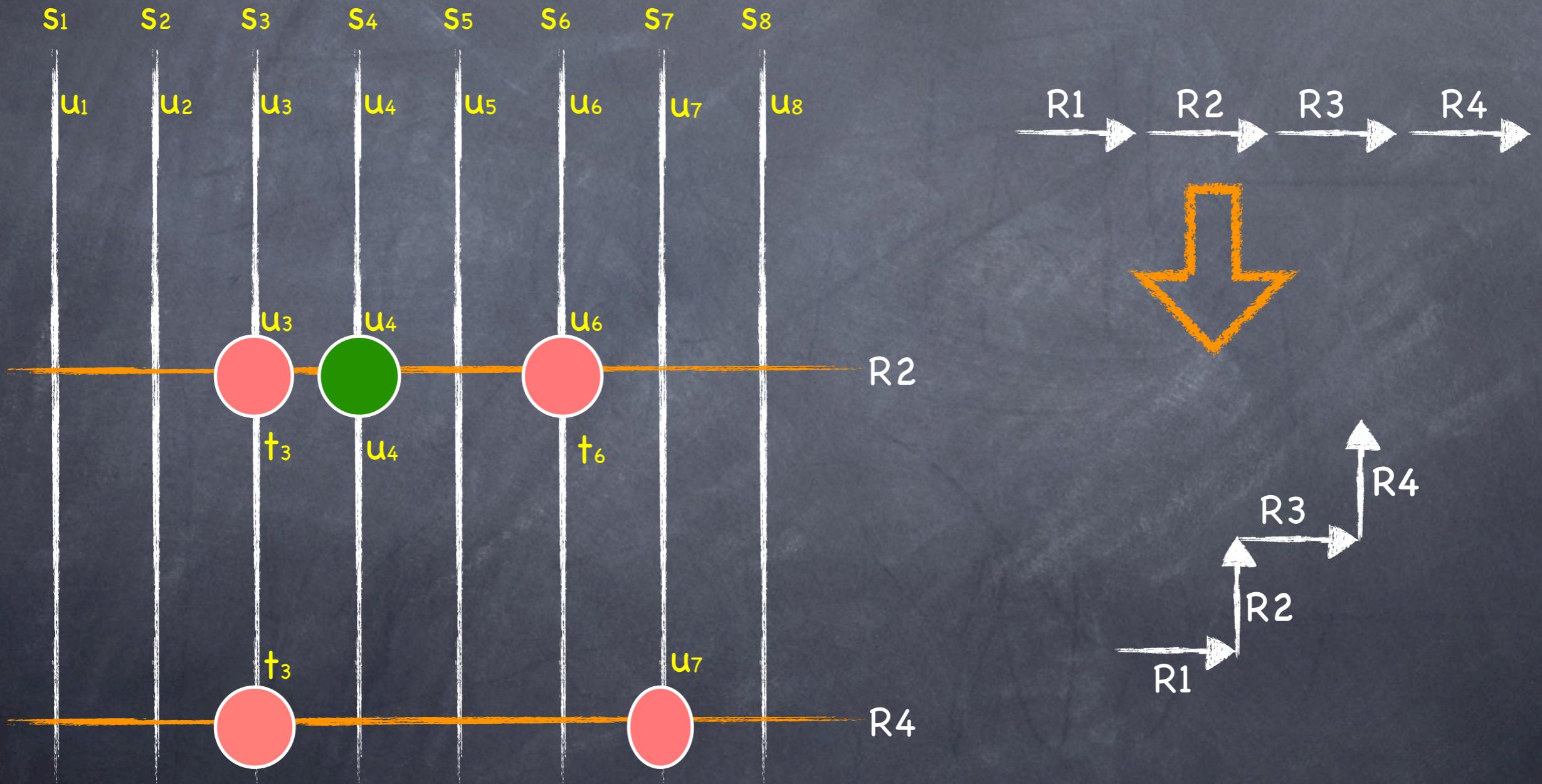
- Concurrency (f is the sequence of all applied rules concurrent to the observable)
- Loop elimination (f is a subset of concurrent rules)
- Helices elimination (a general -action- map)

Say a trace is **hyper causal** if it cannot be compressed further

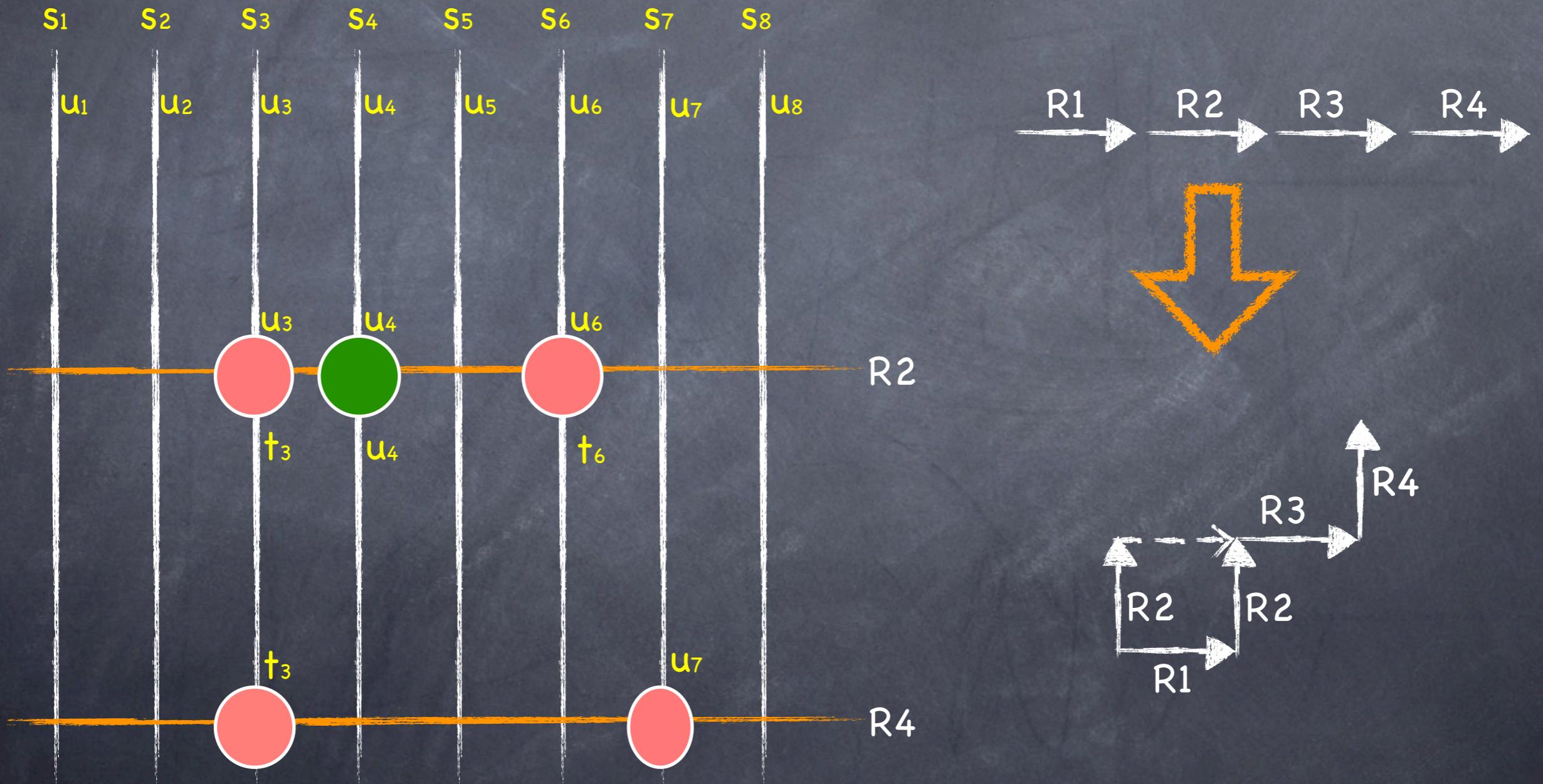
The oracle



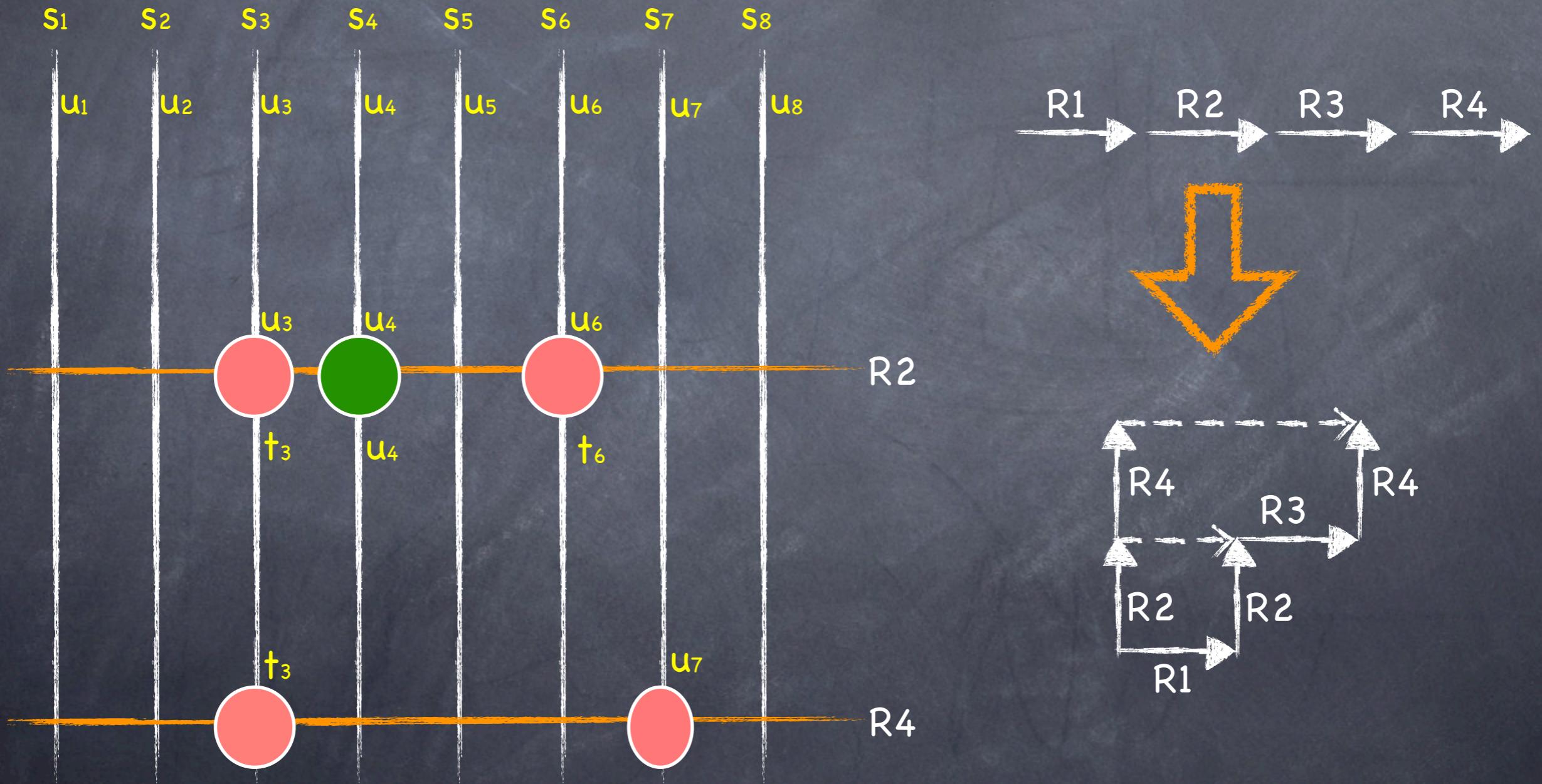
The oracle



The oracle

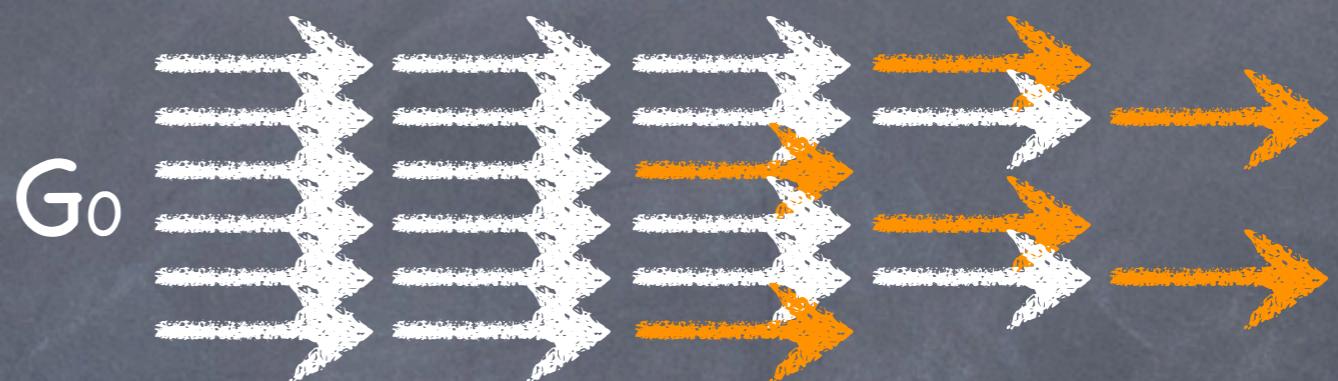


The oracle



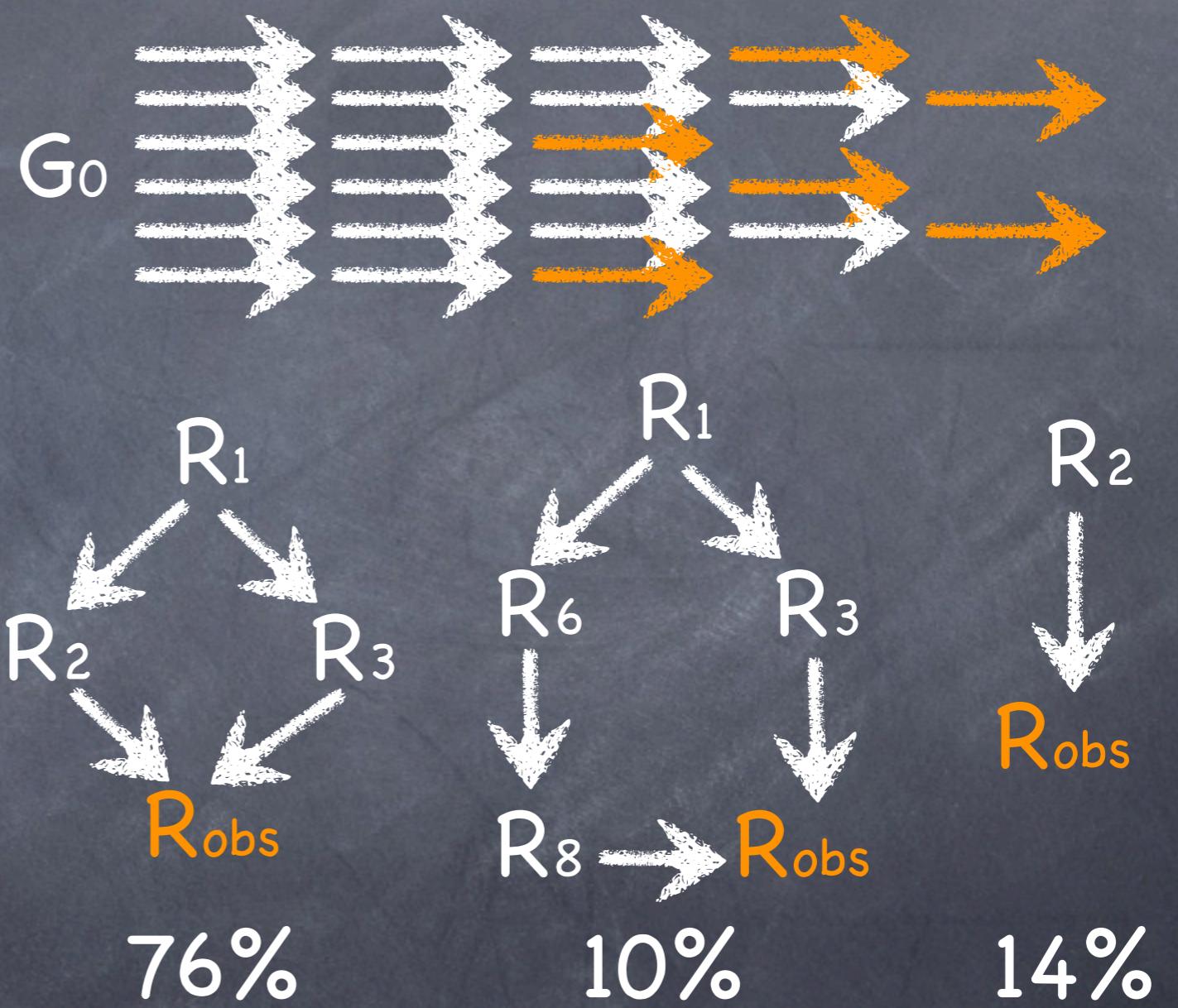
Workflow

- Run n simulations



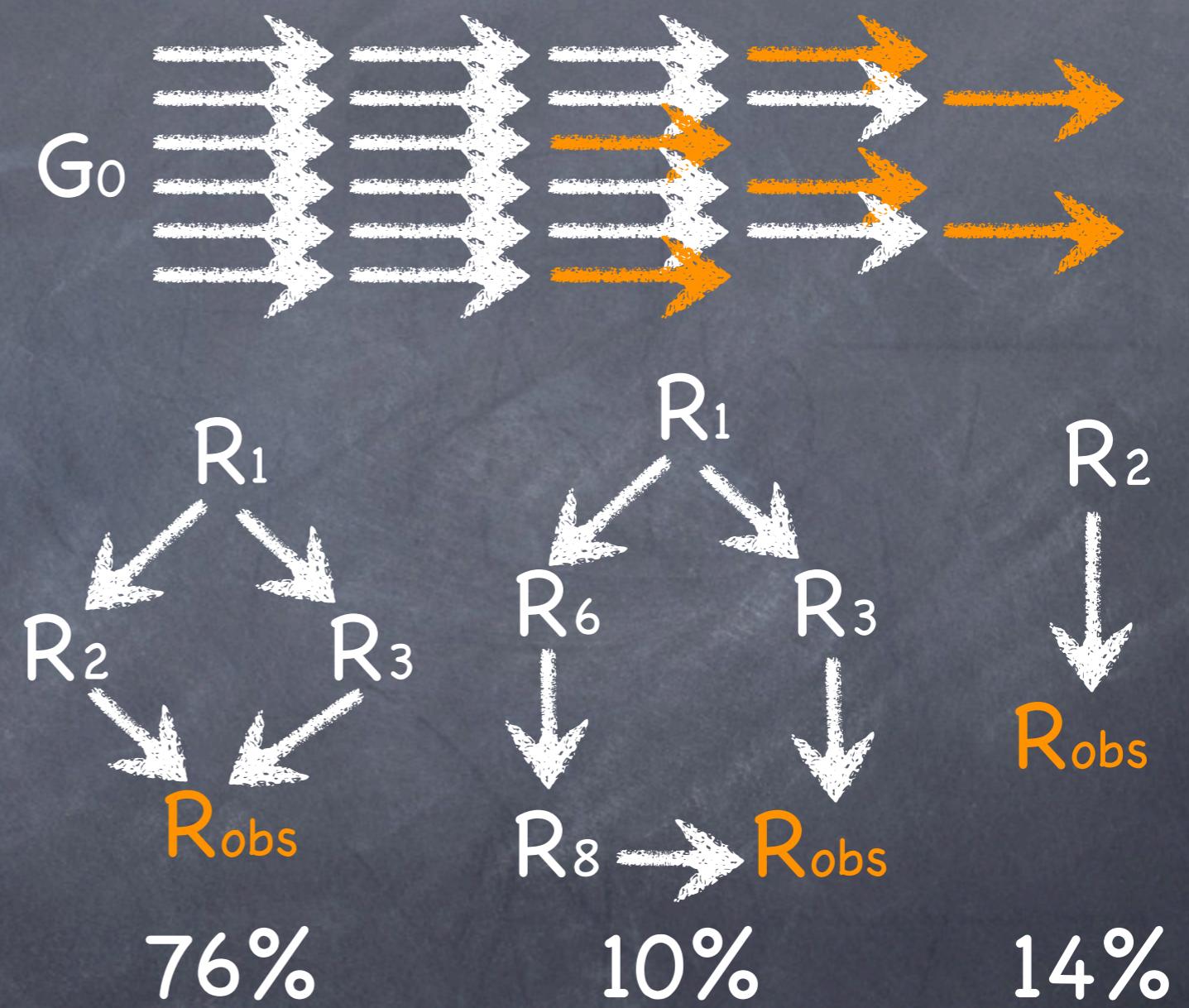
Workflow

- Run n simulations
- Simple causality analysis



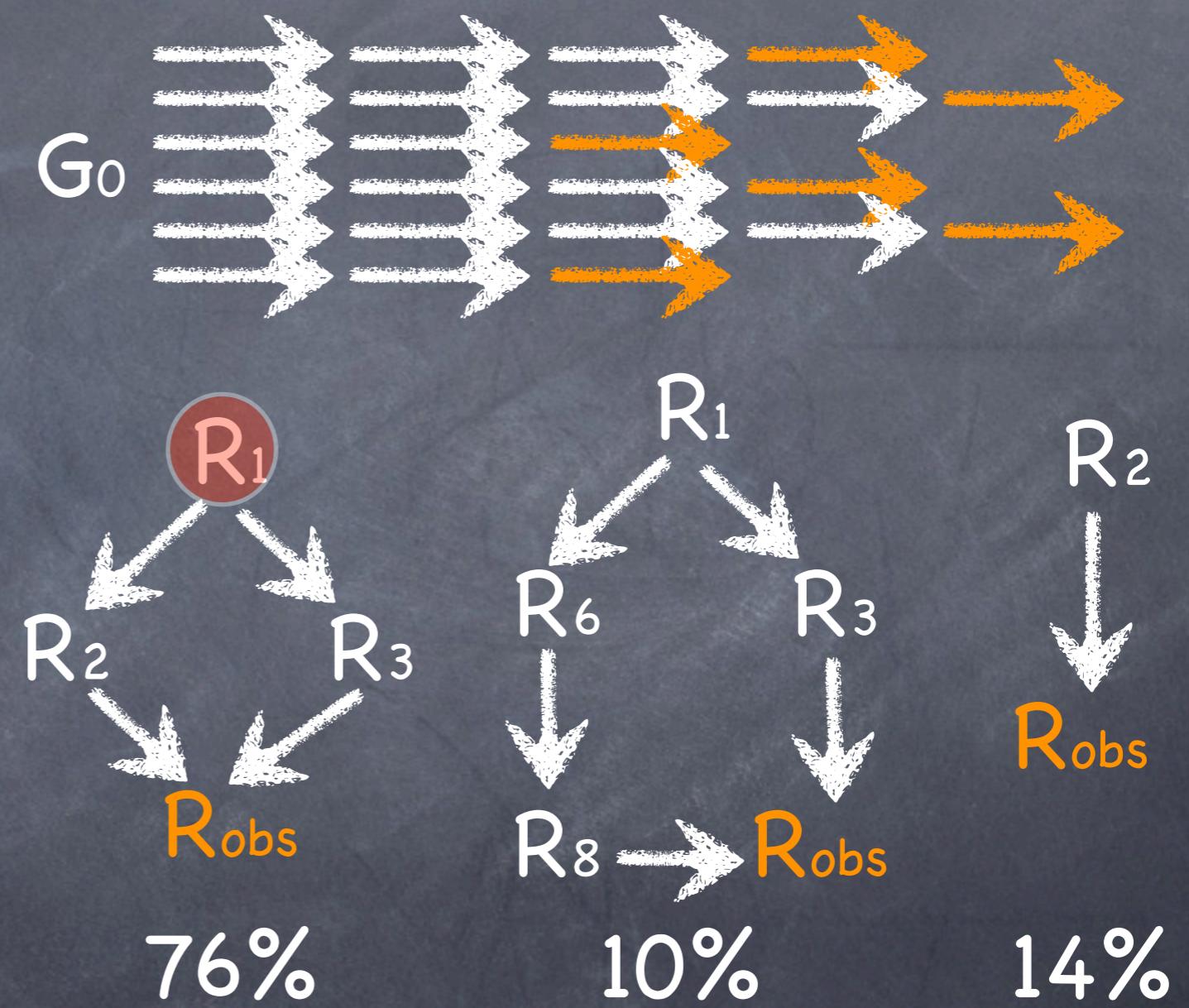
Workflow

- Run n simulations
- Simple causality analysis
- Knock-out events to reduce to minimal configurations



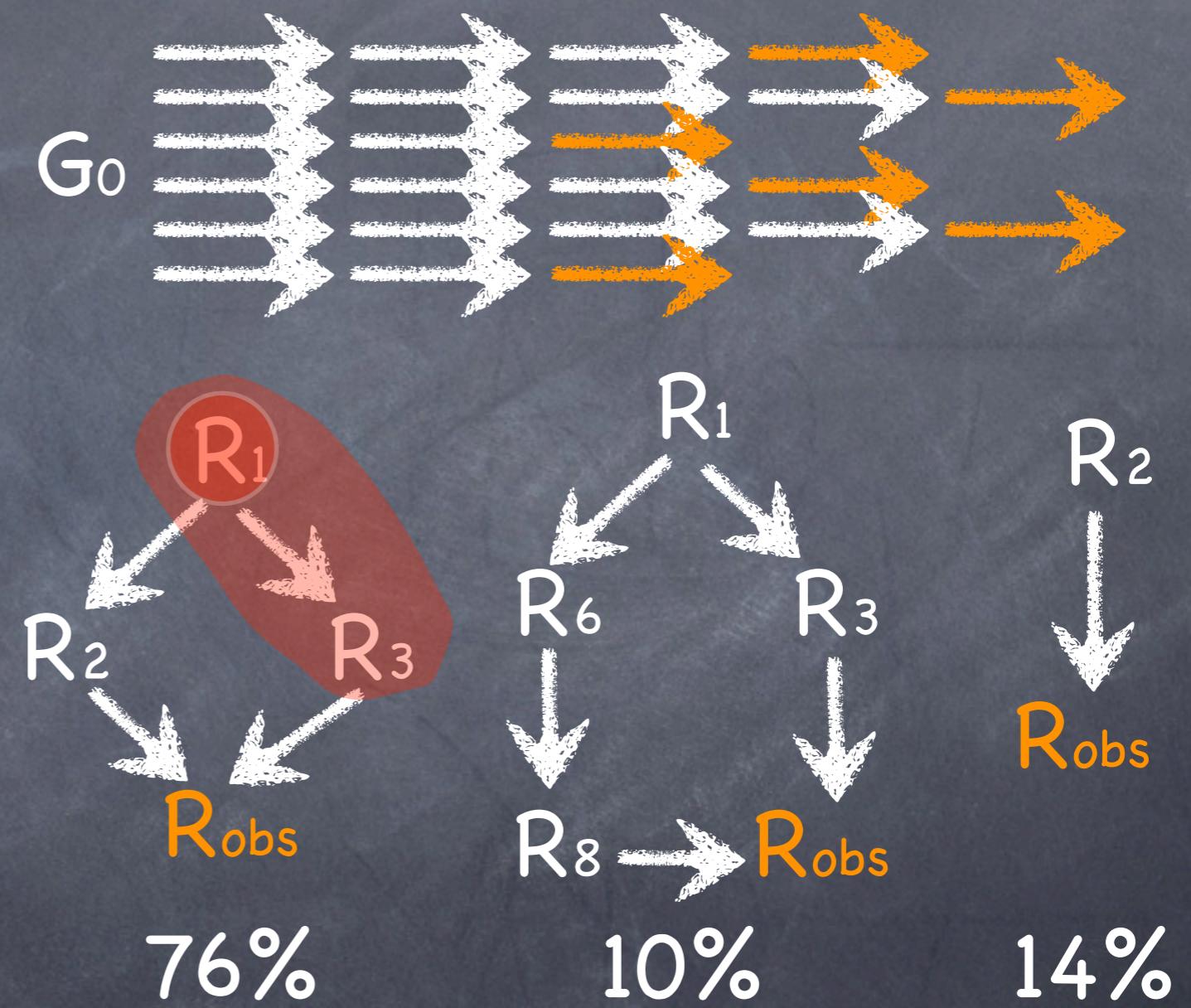
Workflow

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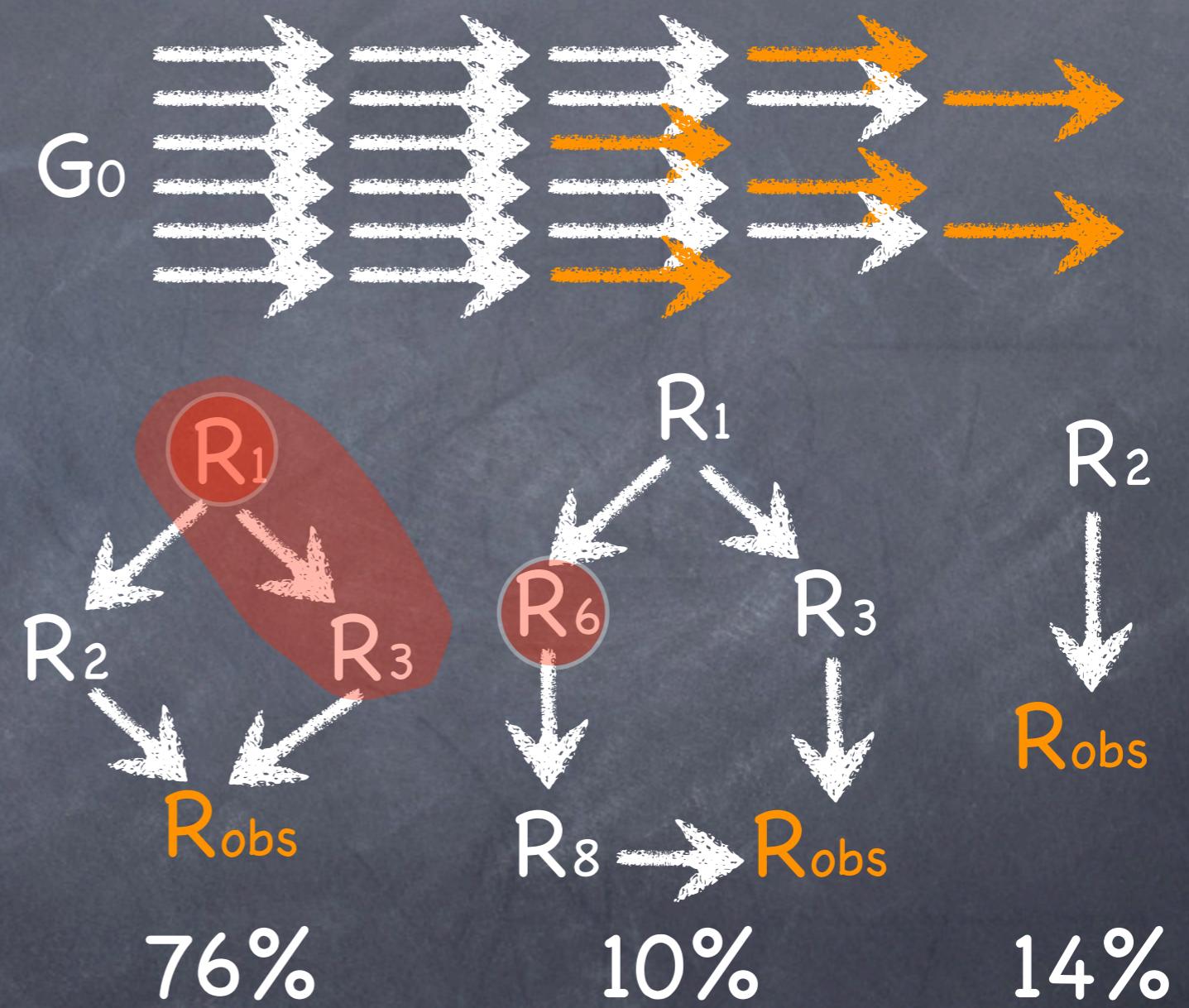
Workflow

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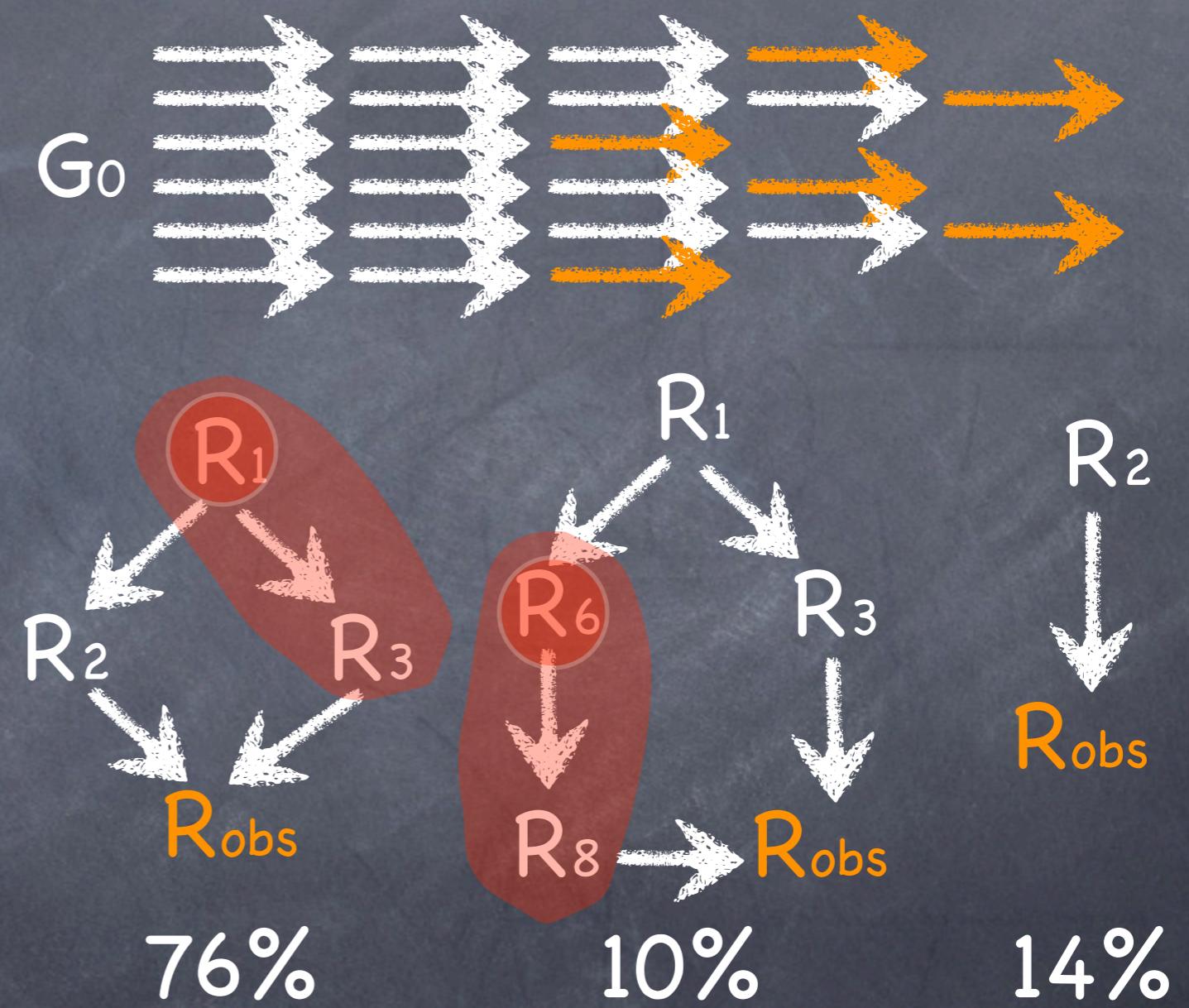
Workflow

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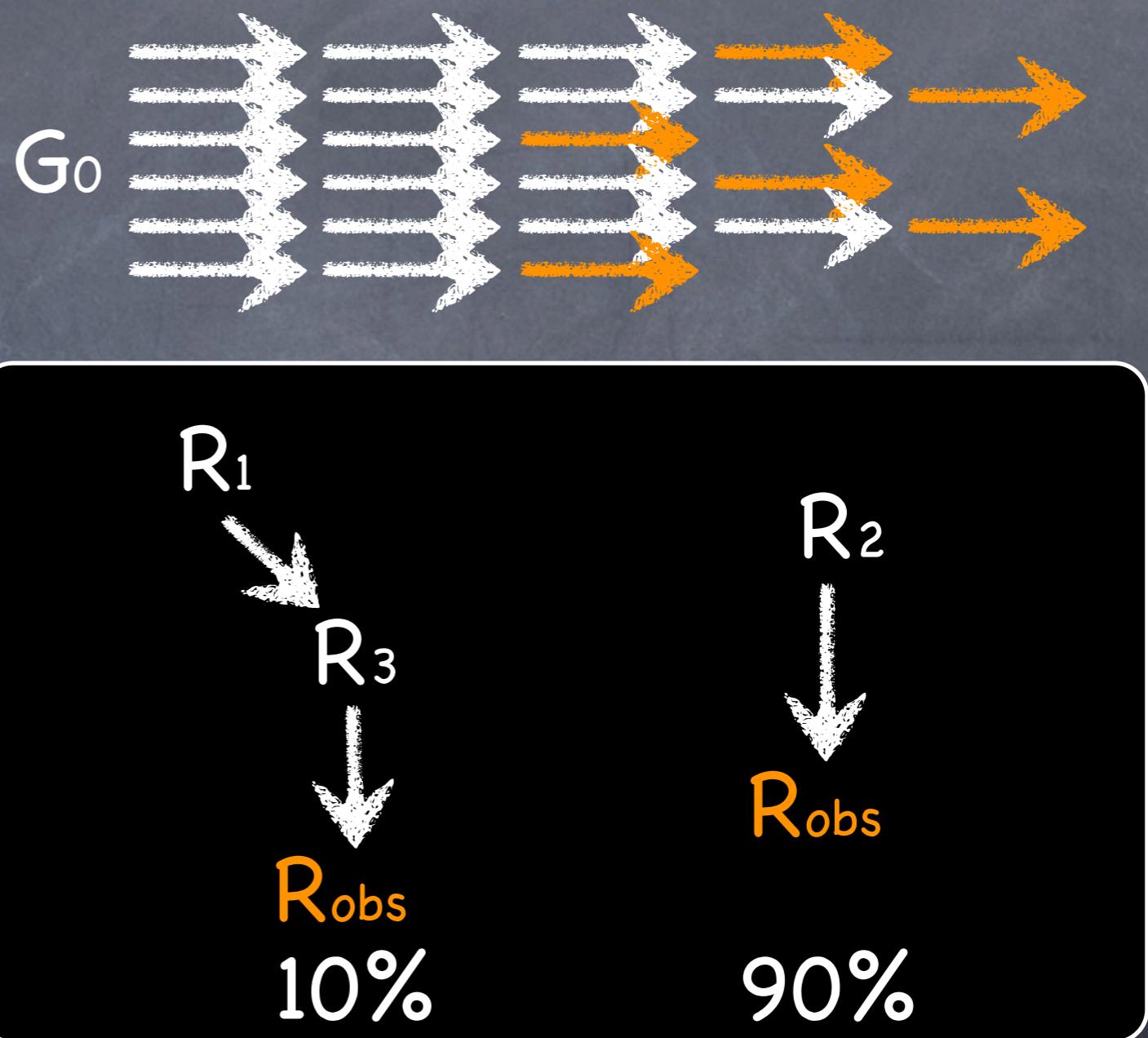
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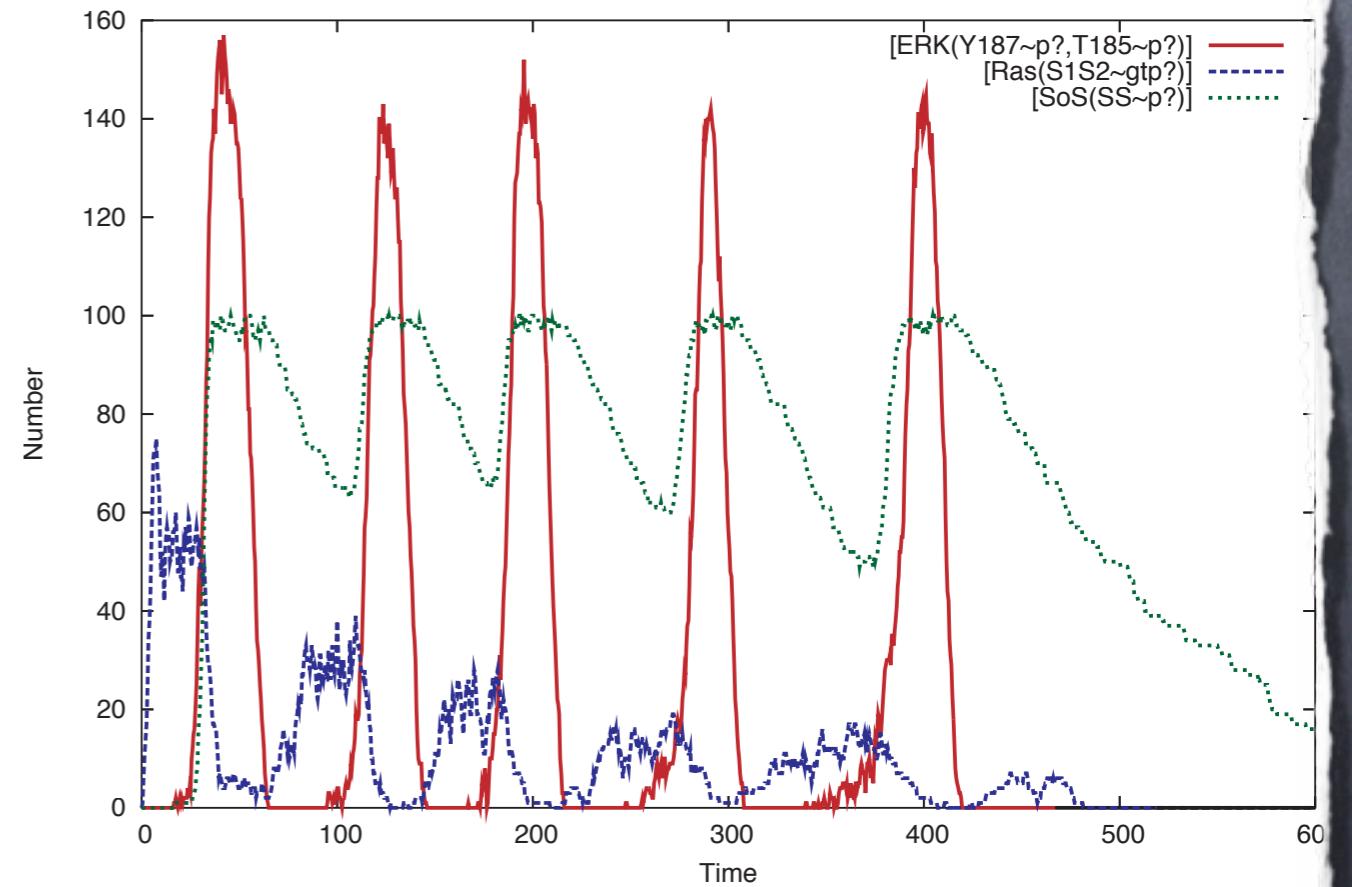
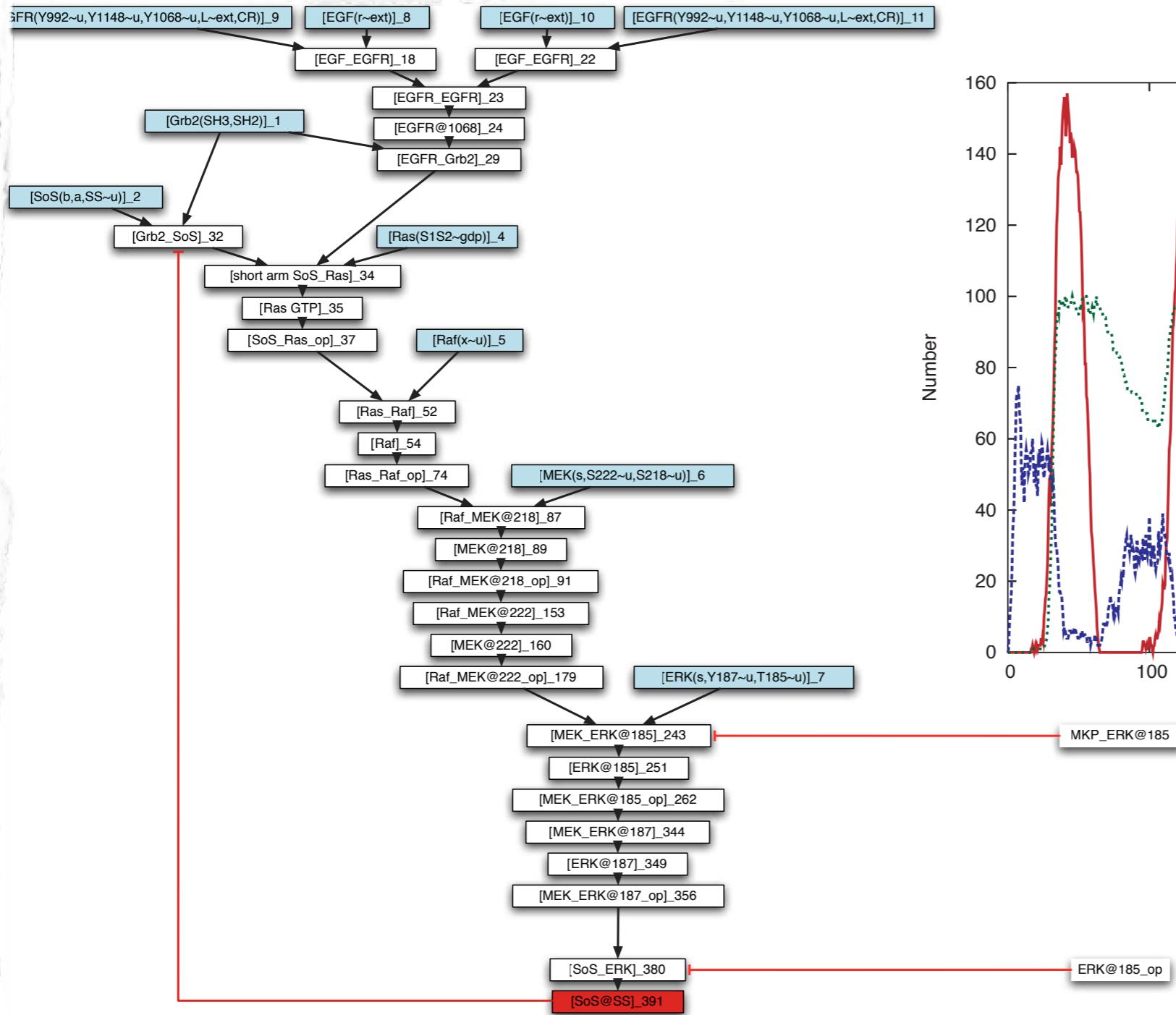


Workflow

- Run n simulations
- Simple causality analysis
- Knock-out events to reduce to minimal configurations



Overview



Temporary conclusions

- ⦿ While exporting standard causality analysis, one discovers theory needs extension!
- ⦿ Tools already useful for debugging purpose
- ⦿ It remains to see whether one can use them for prediction (emergent behavior)

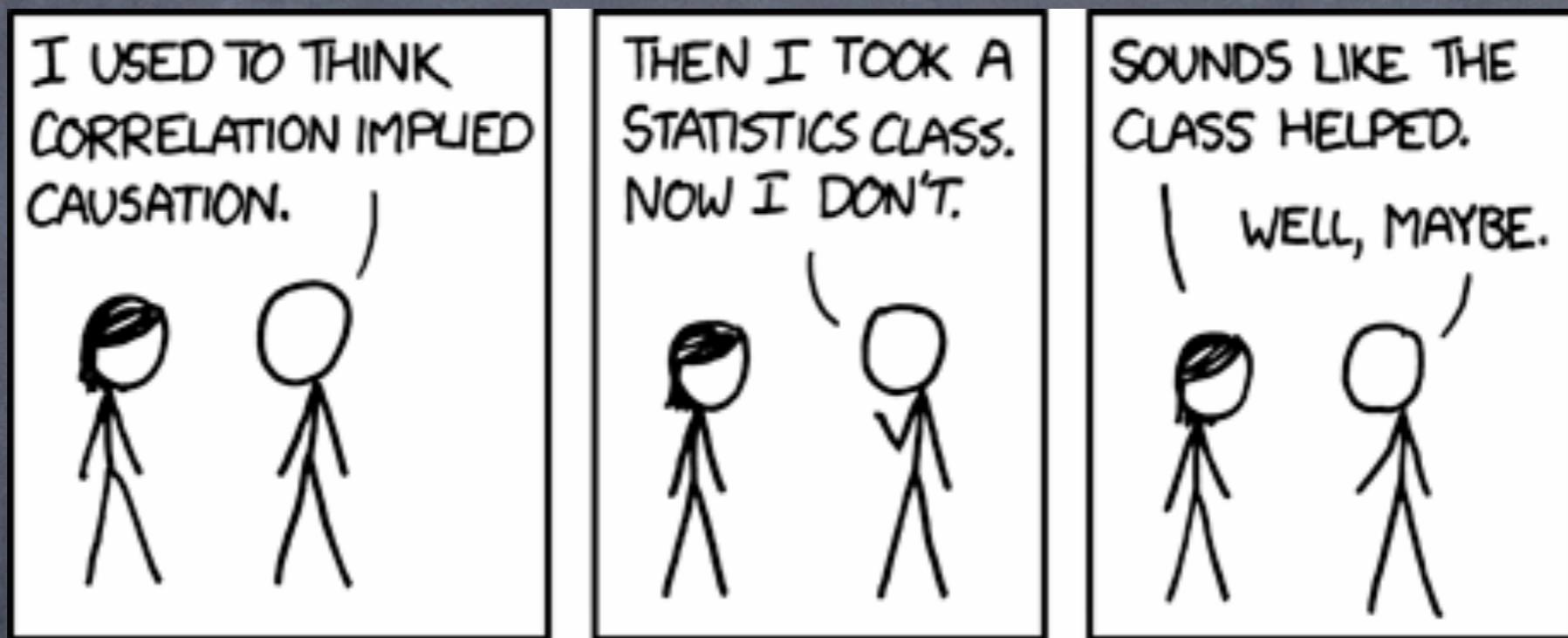
Knock-out in CS

T1	T2	T3	T4
i++	i++	i*i	sqrt i

i:=0
...
i:=nan

Knock out an event and check
whether **obs:nan** is preserved!

Thanks



<http://xkcd.com/>